Porter stemmer

The Porter stemming algorithm is a process for removing suffixes from words in English. Removing suffixes automatically is an operation which is especially useful in the field of information retrieval. In a typical IR environment a document is represented by a vector of words, or terms. Terms with a common stem will usually have similar meanings, for example:

- CONNECT
- CONNECTED
- CONNECTING
- CONNECTION
- CONNECTIONS

Frequently, the performance of an IR system will be improved if term groups such as this are conflated into a single term. This may be done by removal of the various suffixes -ED, -ING, -ION, IONS to leave the single term CONNECT. In addition, the suffix stripping process will reduce the total number of terms in the IR system, and hence reduce the size and complexity of the data in the system, which is always advantageous.

Usually is desired that only inflectional morphemes are removed (those corresponding to declinations, conjugations, etc), not also derivational morphemes (which correspond to different parts of speech). Porter algorithm does not fulfill this. One can make his own set of rules (for any language). Porter’s stemmer advantage is its simplicity and speed.

Porter algorithm was made in the assumption that we don’t have a stem dictionary and that the purpose of the task is to improve IR performance (not as a linguistic exercise). The program is given an explicit list of suffixes, and, with each suffix, the criterion under which it may be removed from a word to leave a valid stem.

Perhaps the best criterion for removing suffixes from two words W1 and W2 to produce a single stem S, is to say that we do so if there appears to be no difference between the two statements ‘a document is about W1’ and ‘a document is about W2’. So if W1=’CONNECTION’ and W2=’CONNECTIONS’ it seems very reasonable to conflate them to a single stem. But if W1=’RELATE’ and W2=’RELATIVITY’ it seems perhaps unreasonable, especially if the document collection is concerned with theoretical physics. Between these two extremes there is a continuum of different cases, and given two terms W1 and W2, there will be some variation in opinion as to whether they should be conflated, just as there is with deciding the relevance of some document to a query. The evaluation of the worth of a suffix stripping system is correspondingly difficult.

The success rate for the suffix stripping will be significantly less than 100% irrespective of how the process is evaluated. For example, if SAND and SANDER get conflated, so most probably will WAND and WANDER. The error here is that the -ER of WANDER has been treated as a suffix when in fact it is part of the stem. Equally, a suffix may completely alter the meaning of a word, in which case its removal is unhelpful. PROBE and PROBATE for example, have quite distinct meanings in modern English. The addition of more rules to increase the performance in one area of the vocabulary causes an equal degradation of performance elsewhere. It is also easy to give undue emphasis to cases which appear to be important, but which turn out to be rather rare. For example, cases in which the root of a word changes with the addition of a suffix, as in DECEIVE/DECEPTION, RESUME/RESUMPTION, INDEX/INDICES occur much more rarely in real vocabularies than one might at first suppose. In view of the error rate that must in any case be expected, it did not seem worthwhile to try and cope with these cases.
THE ALGORITHM

A consonant in a word is a letter other than A, E, I, O or U, and other than Y preceded by a consonant. So in TOY the consonants are T and Y, and in SYZYGY they are S, Z and G. If a letter is not a consonant it is a vowel.

A consonant will be denoted by c, a vowel by v. A list ccc... of length greater than 0 will be denoted by C, and a list vvv... of length greater than 0 will be denoted by V. Any word, or part of a word, therefore has one of the four forms:

CVCV ... C
CVCV ... V
VCVC ... C
VCVC ... V

These may all be represented by the single form:

[C]VCVC ... [V]

where the square brackets denote arbitrary presence of their contents.
Using (VC)m to denote VC repeated m times, this may again be written as:

[C](VC){m}[V].

m will be called the measure of any word or word part when represented in this form. The case m = 0 covers the null word. Here are some examples:

m=0  TR, EE, TREE, Y, BY.
m=1  TROUBLE, OATS, TREES, IVY.
m=2  TROUBLES, PRIVATE, OATEN, ORREY.

The rules for removing a suffix will be given in the form:

(condition) S1 \rightarrow S2

This means that if a word ends with the suffix S1, and the stem before S1 satisfies the given condition, S1 is replaced by S2. The condition is usually given in terms of m, e.g.:

(m > 1) EMENT \rightarrow

Here S1 is ‘EMENT’ and S2 is null. This would map REPLACEMENT to REPLAC, since REPLAC is a word part for which m = 2.

The ‘condition’ part may also contain the following:

*S  - the stem ends with S (and similarly for the other letters).
*V* - the stem contains a vowel.
*d  - the stem ends with a double consonant (e.g. -TT, -SS).

*o  - the stem ends cvc, where the second c is not W, X or Y (e.g. -WIL, -HOP).

And the condition part may also contain expressions with and, or and not, so that:

(m>1 and (*S or *T))

tests for a stem with m > 1 ending in S or T, while

(*d and not (*L or *S or *Z))

tests for a stem ending with a double consonant other than L, S or Z.

Elaborate conditions like this are required only rarely.

In a set of rules written beneath each other, only one is obeyed, and this will be the one with the longest matching S1 for the given word. For example, with:

SSES  ->  SS
IES  ->  I
SS  ->  SS
S  ->

(here the conditions are all null) CARESES maps to CARESS since SSES is the longest match for S1. Equally CARESS maps to CARESS (S1=“SS”) and CARES to CARE (S1=“S”).

In the rules below, examples of their application, successful or otherwise, are given on the right in lower case. The algorithm now follows:

Step 1a

SSES  ->  SS
IES  ->  I
SS  ->  SS
S  ->

caresses  ->  caress
ponies  ->  poni
ties  ->  ti
caress  ->  caress
cats  ->  cat

Step 1b

(m>0) EED  ->  EE

feed  ->  feed
agreed  ->  agree
plastered  ->  plaster
bled  ->  bled
motoring  ->  motor
sing  ->  sing

If the second or third of the rules in Step 1b is successful, the following is done:
AT -> ATE  conflat(ed) -> conflate
BL -> BLE  troubl(ed) -> trouble
IZ -> IZE  siz(ed) -> size
(*d and not (*L or *S or *Z))
  -> single letter
  hopp(ing) -> hop
tann(ed) -> tan
fall(ing) -> fall
hiss(ing) -> hiss
fizz(ed) -> fizz
(m=1 and *o) -> E
  fail(ing) -> fail
  fil(ing) -> file

The rule to map to a single letter causes the removal of one of the double letter pair. The -E is put back on -AT, -BL and -IZ, so that the suffixes -ATE, -BLE and -IZE can be recognized later. This E may be removed in step 4.

Step 1c

(*v*) Y -> I
  happy -> happi
  sky -> sky

Step 1 deals with plurals and past participles. The subsequent steps are much more straightforward.

Step 2

(m>0) ATIONAL -> ATE  relational -> relate
(m>0) TIONAL -> TION  conditional -> condition
                rational -> rational
(m>0) ENCI -> ENCE  valenci -> valence
(m>0) ANCI -> ANCE  hesitanci -> hesitance
(m>0) IZER -> IZE  digitizer -> digitize
(m>0) ABLI -> ABLE  conformabli -> conformable
(m>0) ALLI -> AL  radicalli -> radical
(m>0) ENTLI -> ENT  differentli -> different
(m>0) ELI -> E  vileli -> vile
(m>0) OUSLI -> OUS  analogousli -> analogous
(m>0) IZATION -> IZE  vietnamization -> vietnamize
(m>0) ATION -> ATE  predication -> predicate
(m>0) ATOR -> ATE  operator -> operate
(m>0) ALISM -> AL  feudalism -> feudal
(m>0) IVENESS -> IVE  decisiveness -> decisive
(m>0) FULNESS -> FUL  hopefulness -> hopeful
(m>0) OUSNESS -> OUS  callousness -> callous
(m>0) ALITI -> AL  formaliti -> formal
(m>0) IVITI -> IVE  sensitiviti -> sensitive
(m>0) BILITI -> BLE  sensibiliti -> sensible

The test for the string S1 can be made fast by doing a program switch on the penultimate letter of the word being tested. This gives a fairly even breakdown of the possible values of the string S1. It will be seen in fact that the S1-strings in step 2 are presented here in the alphabetical order of their penultimate letter.
Similar techniques may be applied in the other steps.

Step 3

\( m > 0 \)

- ICATE -> IC  ttriplicate -> triplic
- ATIVE -> form
- ALIZE -> AL formalize -> formal
- ICITI -> IC electriciti -> electric
- ICAL -> IC electrical -> electric
- FUL -> hopeful -> hope
- NESS -> goodness -> good

Step 4

\( m > 1 \)

- AL -> revival -> reviv
- ANCE -> allowance -> allow
- ENCE -> inference -> infer
- ER -> airliner -> airlin
- IC -> gyrosopic -> gyroscope
- ABLE -> adjustable -> adjust
- IBLE -> defensible -> defens
- ANT -> irritant -> irrit
- EMENT -> replacement -> replac
- MENT -> adjustment -> adjust
- ENT -> dependent -> depend
- ION -> adoption -> adopt
- ISM -> communism -> commun
- ATE -> activate -> activ
- ITI -> angulariti -> angular
- OUS -> homologous -> homolog
- IVE -> effective -> effect
- IZE -> bowdlerize -> bowdler

The suffixes are now removed. All that remains is a little tidying up.

Step 5a

\( m > 1 \)

- E -> probate -> probat
- 0U -> rate -> rate

\( m = 1 \) and not *o* E ->

- cease -> ceas

Step 5b

\( m > 1 \) and *d and *L -> single letter

- controll -> control
- roll -> roll

The algorithm is careful not to remove a suffix when the stem is too short, the length of the stem being given by its measure, \( m \). There is no linguistic basis for this approach. It was merely observed that \( m \) could be used quite effectively to help decide whether or not it was wise to take off a suffix.
For example, in the following two lists:

<table>
<thead>
<tr>
<th>list A</th>
<th>list B</th>
</tr>
</thead>
<tbody>
<tr>
<td>------</td>
<td>------</td>
</tr>
<tr>
<td>RELATE</td>
<td>DERIVATE</td>
</tr>
<tr>
<td>PROBATE</td>
<td>ACTIVATE</td>
</tr>
<tr>
<td>CONFLATE</td>
<td>DEMONSTRATE</td>
</tr>
<tr>
<td>PIRATE</td>
<td>NECESSITATE</td>
</tr>
<tr>
<td>PRELATE</td>
<td>RENOVATE</td>
</tr>
</tbody>
</table>

-ATE is removed from the list B words, but not from the list A words. This means that the pairs DERIVATE/DERIVE, ACTIVATE/ACTIVE, DEMONSTRATE/Demonstrate, NECESSITATE/NECESSITOUS, will conflate together. The fact that no attempt is made to identify prefixes can make the results look rather inconsistent. Thus PRELATE does not lose the -ATE, but ARCHPRELATE becomes ARCHPREL.

Complex suffixes are removed bit by bit in the different steps. Thus GENERALIZATIONS is stripped to GENERALIZATION (Step 1), then to GENERALIZE (Step 2), then to GENERAL (Step 3), and then to GENER (Step 4). OSCILLATORS is stripped to OSCILLATOR (Step 1), then to OSCILLATE (Step 2), then to OSCILL (Step 4), and then to OSCIL (Step 5). In a vocabulary of 10,000 words, the reduction in size of the stem was distributed among the steps as follows:

<table>
<thead>
<tr>
<th>Suffix stripping of a vocabulary of 10,000 words</th>
</tr>
</thead>
<tbody>
<tr>
<td>Number of words reduced in step 1: 3597</td>
</tr>
<tr>
<td>&quot; 2: 766</td>
</tr>
<tr>
<td>&quot; 3: 327</td>
</tr>
<tr>
<td>&quot; 4: 2424</td>
</tr>
<tr>
<td>&quot; 5: 1373</td>
</tr>
<tr>
<td>Number of words not reduced: 3650</td>
</tr>
</tbody>
</table>

The resulting vocabulary of stems contained 6370 distinct entries. Thus the suffix stripping process reduced the size of the vocabulary by about one third.

Similar algorithms can be applied for different languages, one can define his own rules adapted to his own needs.