

NOT NULL ICs:

- Example:

```
CREATE TABLE Drinkers(  
    name CHAR(30) PRIMARY KEY,  
    phone CHAR(16) NOT NULL );
```

- Otherwise, without the **NOT NULL** IC, we get:

```
INSERT INTO Drinkers(name)  
VALUES('Sally');
```

<u>name</u>	<u>phone</u>
Sally	NULL

- Other specifications:
 - **NOT NULL:** every tuple must have a real value for this attribute
 - **DEFAULT value:** a value to use whenever no other value for this attribute is known
- Example:

```
CREATE TABLE Drinkers (  
    name CHAR(30) PRIMARY KEY,  
    addr CHAR(50) DEFAULT '123 Sesame St',  
    phone CHAR(16) );
```

- With the insertion:

```
INSERT INTO Drinkers(name)
VALUES('Sally');
```

name	addr	phone
Sally	123 Sesame St.	NULL

- Primary key is by default NOT NULL
- This insertion is perfectly legal
It is OK to list a subset of the attributes, and values for only this subset
- Had we declared
 “`phone CHAR(16) NOT NULL`”,
the insertion could not be done

Changing Columns of a Schema:

- Add an attribute to relation R with

```
ALTER TABLE R ADD <column declaration>;
```

- Example:

```
ALTER TABLE Bars ADD phone CHAR(16)  
DEFAULT 'unlisted';
```

- Columns may also be dropped:

```
ALTER TABLE Bars DROP license;
```

SQL: Queries

- Basic Syntax:

SELECT desired attributes
FROM tuple variables - range over relations
WHERE condition about tuple variables

- Example: Schema:

Beers(name, manf)
Bars(name, addr, license)
Drinkers(name, addr, phone)
Likes(drinker, beer)
Sells(bar, beer, price)
Frequents(drinker, bar)

- Query: What beers are made by Anheuser-Busch?

```
SELECT name
FROM Beers
WHERE manf = 'Anheuser-Busch';
```

(notice the single quotes for strings)

- Answer:

name
Bud
Bud Lite
Michelob

- Conditions in WHERE capture RA selection

SELECTing only some attributes captures RA projection

- Extensions of the Basic Syntax:
- A star can be used to retrieve all the attributes:

Beers(name, manf)

```
SELECT *
FROM Beers
WHERE manf = 'Anheuser-Busch';
```

- Answer:

name	manf
Bud	Anheuser-Busch
Bud Lite	Anheuser-Busch
Michelob	Anheuser-Busch

- Renaming Columns: 'name' into 'beer' in Beers(name, manf)

```
SELECT name AS beer
FROM Beers
WHERE manf = 'Anheuser-Busch';
```

- Answer:

beer
 Bud
 Bud Lite
 Michelob

Table useful by itself

Can be combined with other queries

- Expressions as Values in Columns:

- Sells(bar, beer, price)

```
SELECT bar, beer, price*150 AS priceInYen
FROM Sells;
```

'price*150': expression 'priceInYen': name for it

- Answer:

bar	beer	priceInYen
Joe's	Bud	300
Sue's	Miller	360

- Answers with a particular string in each row: use that constant as an expression Likes(drinker, beer)

```
SELECT drinker, 'likes Bud' AS whoLikesBud
FROM Likes
WHERE beer = 'Bud';
```

'likes Bud': expression 'whoLikesBud': name for it

- Answer:

drinker	whoLikesBud
Sally	likes Bud
Fred	likes Bud

- Example: Find the price Joe's Bar charges for Bud

- `Sells(bar, beer, price)`

```
SELECT price
FROM Sells
WHERE bar = 'Joe''s Bar' AND beer = 'Bud';
```

- Two single-quotes in a string represent a single quote
- Conditions in WHERE clause can use logical connectives AND, OR, NOT and parentheses as usual
- SQL is case insensitive
Keywords like SELECT or AND in upper or lower case
Only inside quoted strings does case matter

- Patterns:
- '%' stands for any string
- '_' stands for any (single) character
- "Attribute **LIKE** pattern" is a condition that is true if the string value of the attribute matches the pattern
Also **NOT LIKE** can be used
- Example: Find drinkers whose phone starts with 555

```
Drinkers(name, addr, phone)
```

```
SELECT name  
FROM Drinkers  
WHERE phone LIKE '%555-____';
```

- Patterns must be quoted, like strings

- Multi-Relation Queries:

- A list of relations in FROM clause

- Notation: *Relation.Attribute*

Disambiguates shared attributes from several relations

- Example: Find the beers that the frequenters of Joe's Bar like

`Likes(drinker, beer)` `Frequents(drinker, bar)`

'drinker' is in two different tables

```
SELECT beer
FROM   Frequents, Likes
WHERE  bar = 'Joe''s Bar' AND
       Frequents.drinker = Likes.drinker;
```

- We can express products and joints of RA
- Here: Likes \bowtie Frequents, selection, projection (on beer)
- The common class of “select-project-join” (SPJ) queries
Also known as “conjunctive queries”

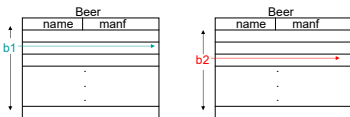
- Explicit Tuple Variables:
- Sometimes we need to refer to two or more copies of a relation
E.g. for comparing values within a same relation
- Use explicit variables for copies of relations (or tuples thereof)
Intuitively: Use *tuple variables* as aliases for the relations
- Example: Find pairs of beers by the same manufacturer

Beers(name, manf)

```

SELECT b1.name, b2.name
→ FROM Beers b1, Beers b2
WHERE b1.manf = b2.manf AND
      b1.name <> b2.name;

```



- **b1, b2**: tuple variables, aliases for relations b1.manf = b2.manf ?

In FROM clause: selection from copies **b1** and **b2** of Beers

- '**b1.name <> b2.name**' needed to avoid producing (Bud, Bud), and a same pair in both orders

- Before we used AS to rename attributes or expressions
- SQL permits AS between a relation and its tuple variable:

```
SELECT b1.name, b2.name
FROM   Beers AS b1, Beers AS b2
WHERE  b1.manf = b2.manf AND b1.name <> b2.name;
```

- Same query in RC:

$$Ans(x, u): \exists y(Beer(x, y) \wedge Beer(u, y) \wedge x \neq u)$$

- A conjunctive query in RC, with a built-in (\neq)
- Exercise: Pose the query in RA

- Sub-Queries:
- Result of a *select-from-where* query can be used in the WHERE clause of another query

The Simplest Case: Sub-query returns a single tuple

- Example: Find bars that serve Miller at the same price Joe charges for Bud

Sells(bar, beer, price)

- First we find what Joe charges for Bud
Next, the bars that sell Miller for that same price

```
SELECT bar
FROM Sells
WHERE beer = 'Miller' AND price =
      (SELECT price
       FROM Sells
       WHERE bar = 'Joe''s Bar' AND beer = 'Bud'
      );
```

- Nested SELECTs

```
SELECT bar
FROM Sells
WHERE beer = 'Miller' AND price =
      (SELECT price
       FROM Sells
       WHERE bar = 'Joe''s Bar' AND beer = 'Bud'
      );
```

- *Scoping rule*: An attribute refers to the closest nested relation with that attribute
- Parentheses around sub-queries are essential
- `price` is uniquely determined (by the key)
The sub-query returns a single value

- What if the sub-query returns a set of tuples?
- Use “**IN**” instead of ‘=’

It is true iff the tuple is in the extension of the relation defined by sub-query

- Example: Find the name and manufacturer of beers that Fred likes

Beers(name, manf)

Likes(drinker, beer)

```
SELECT *
FROM Beers
WHERE name IN
      (SELECT beer
       FROM Likes
       WHERE drinker = 'Fred');
```

- Main (final) selection is on **Beers**, with attributes (**name,manf**)
But the condition depends on table **Likes** (in the sub-query)
- The sub-query creates a set (a unary relation of names)
Main query checks if **name** belongs to it
- In general (not here), **NOT IN** can also be used

- Quantifiers:

- Inherited from Relational Calculus

“**EXISTS** (Relation)” is true iff Relation is non-empty

- Example: Find the beers that are the only beer made by its company

Beers(name, manf)

```
SELECT b1.name
FROM   Beers b1
WHERE  NOT EXISTS(
      SELECT *
      FROM   Beers
      WHERE  manf = b1.manf
            AND name <> b1.name );
```

“Choose the name of the beers such that there is no other beer produced by the same manufacturer”

- Think of having a generic, variable, tuple in **Beers**, **b1**, and checking a condition on it

A condition in terms of a quantifiers (**NOT EXISTS**) applied to a sub-query

Keep the **b1**s that satisfy the condition (their **names**)

```

SELECT b1.name
FROM Beers b1
WHERE NOT EXISTS(
  SELECT *
  FROM Beers
  WHERE manf = b1.manf
        AND name <> b1.name );

```

- Variable `b1` goes over tuples in `Beers` (alias for `Beers`)
- `b1.name` is just another name to refer to `name`
- In the subquery, with `b1` fixed (as a variable outside the subquery):
We get the tuples (`name`, `manf`) with:
$$\text{manf} = \text{b1.manf} \ \& \ \text{name} \neq \text{b1.name}$$
- *Scoping rule*: To refer to external `Beers` in the sub-query, give to the external tuple a variable (`b1` in this example)
- A subquery that refers to values from a surrounding query is called a *correlated sub-query*
- Exercise: Express the query in RC and RA

- **ANY** and **ALL** behave as **existential and universal quantifiers**, respectively
- Beware: In English, “any” and “all” sometimes act as synonyms
 For example, “I am fatter than any of you” vs. “I am fatter than all of you” **Not in SQL!**

- They can be used to express numerical maxima and minima
 “A value is a maximum if *all* values are not higher”

- Example: Find the beers sold for the highest price

```

                Sells(bar, beer, price)
SELECT beer
FROM   Sells
WHERE  price >= ALL(SELECT price           (the sub-query returns real numbers)
                   FROM   Sells);

```

- Exercise: Find the beers not sold for the lowest price
 (use ANY)

- Bank Example: Schema:

```
Branch(branch-name, branch-city, assets)
Customer(customer-name, customer-street, customer-city)
Account(branch-name, account#, balance)
Depositor(customer-name, account#)
Loan(branch-name, loan#, amount)
Borrower(customer-name, loan#)
```

- **Example 1:** Find all customers who have both an account and a loan at the bank

```
SELECT customer-name
FROM Borrower
WHERE customer-name IN (SELECT customer-name
                        FROM Depositor);
```

- **Example 2:** Find all customers who have a loan but not an account at the bank

```
SELECT customer-name
FROM Borrower
WHERE customer-name NOT IN (SELECT customer-name
                            FROM Depositor);
```

- **Example 3:** Find the customers at the “First Street” branch with the highest account balance at that branch

```
SELECT customer-name
FROM Depositor, Account
WHERE Depositor.account#=Account.account AND
      branch-name= 'First Street' AND
      Account.balance >= ALL
      (SELECT balance
       FROM Account
       WHERE branch-name='First Street');
```

- Union, Intersection, Difference:
- One can bring RA operators explicitly into SQL
 UNION, INTERSECT as usual
 EXCEPT for the difference of the two relations
- They require a shared schema
- Example: Find the drinkers and beers such that the drinker likes the beer and frequents a bar that serves it

```

Likes(drinker, beer)    Sells(bar, beer, price)
Frequents(drinker, bar)

```

```

(SELECT * FROM Likes)
      INTERSECT
(SELECT drinker, beer
 FROM Sells, Frequents
 WHERE Frequents.bar = Sells.bar );

```

(a way to pick tuples from LIKES
with conditions coming from other relations)
(selection of drinkers that frequent a bar where
those beers are served)

- In SELECT: A join between Sells and Frequents, and a projection over drinker and beer (so, same schema as LIKES)
- Exercise: Solve it without INTERSECT