

THE FORMAL LANGUAGE L_t AND TOPOLOGICAL PRODUCTS

by L. E. BERTOSSI in Santiago (Chile)

1. Preliminaries

We will consider topological structures of the form $\langle \mathfrak{A}, \sigma \rangle$, where \mathfrak{A} is a classical algebraic structure for first order logic and σ is a topology or a basis for a topology on the universe A of \mathfrak{A} . Given a set of symbols S compatible with the algebraic parts of the structures under consideration, we construct the formal language L_t^S for expressing properties of those topological structures. L_t^S is constructed as the classical first-order language L^S adding to it the following rules for building formulas:

if t is a (first-order) term of L^S , then $t \in X$ (where X is a second-order variable for open sets) is an atomic formula;

if φ is a formula of L_t^S positive in X and t is a term of L^S , then $(\forall X \ni t) \varphi$ is a formula of L_t^S ;

if φ is a formula of L_t^S negative in X and t is a term of L^S , then $(\exists X \ni t) \varphi$ is a formula of L_t^S .

It is easy to show that the formulas of L_t^S are basis-invariant, i.e. for each $\varphi \in L_t^S$ and structure $\langle \mathfrak{A}, \sigma \rangle$:

$$\langle \mathfrak{A}, \sigma \rangle \models \varphi \quad \text{iff} \quad \langle \mathfrak{A}, \bar{\sigma} \rangle \models \varphi.$$

Here $\bar{\sigma}$ is the topology generated by σ : $\bar{\sigma} = \{\bigcup s \mid s \subseteq \sigma\}$. For historical remarks, precise definitions and other interesting properties of L_t , see [4], [7].

2. The topological product

Let $\langle \mathfrak{A}_i, \sigma_i \rangle_{i \in I}$ be topological structures and $D \subseteq P(I)$ a filter on I . We define the *D-product of the* $\langle \mathfrak{A}_i, \sigma_i \rangle$ as the topological structure $\langle \prod_I \mathfrak{A}_i, \bar{\sigma}_D \rangle$, where $\prod_I \mathfrak{A}_i$ is the direct product of the \mathfrak{A}_i and $\bar{\sigma}_D$ is the topology generated by the basis

$$\sigma_D = \left\{ \prod_I U_i \mid U_i \in \sigma_i \text{ and } \{i \in I \mid U_i = A_i\} \in D \right\}.$$

Some facts that are easy to verify are the following:

- a) The projection maps p_i are continuous iff $D \supseteq Co(I)$, the filter of cofinite sets.
- b) If $D = Co(I)$, then $\bar{\sigma}_D$ is the Tychonov topology (the usual product topology).
- c) If $D = \{I\}$ (the trivial filter), then $\bar{\sigma}_D$ is the trivial topology.
- d) If $D = P(I)$ (the improper filter), then $\bar{\sigma}_D$ is the box topology.
- e) If D, E are filters on I and $D \subseteq E$, then $\bar{\sigma}_E$ is finer than $\bar{\sigma}_D$.
- f) $\langle \prod_I \mathfrak{A}_i, \bar{\sigma}_D \rangle$ is Hausdorff iff $D \supseteq Co(I)$ and all $\langle \mathfrak{A}_i, \sigma_i \rangle$ are Hausdorff.

If $\mathfrak{B}(I)$ is the Boolean algebra $\langle P(I), \cap, \cup, ^c, \emptyset, I \rangle$, then the structure $\langle \mathfrak{B}(I), D \rangle$ is a Boolean algebra with distinguished filter D . Let L^{BAF} be the first-order language for such algebras.

It is also possible to extend the Feferman-Vaught-type theorem [3] given in [4] for usual topological products to our products:

Theorem. *To each sentence $\varphi \in L_t^S$ there can be effectively associated a tuple $(\Phi; \theta_1, \dots, \theta_n)$, where $\Phi(z_1, \dots, z_n) \in L^{BAF}$ and $\theta_1, \dots, \theta_n \in L_t^S$ are sentences, such that for all $I, (\langle \mathfrak{A}_i, \sigma_i \rangle)_{i \in I}, D$:*

$$\langle \prod_i \mathfrak{A}_i, \bar{\sigma}_D \rangle \models \varphi \text{ iff } \langle \mathfrak{B}(I), D \rangle \models \Phi[S(\theta_1), \dots, S(\theta_n)],$$

where $S(\theta_j) = \{i \in I \mid \langle \mathfrak{A}_i, \sigma_i \rangle \models \theta_j\}$. \square

As in [3], we can everywhere assume that $(\Phi; \theta_1, \dots, \theta_n)$ is a partitioning tuple.

Combining this result with ERSHOV's theorem [2] on the decidability of the first-order theory of all algebras of the form $\langle \mathfrak{B}(I), D \rangle$, we obtain

Theorem. *If K is a class of topological structures with decidable L_t^S -theory, then the L_t^S -theory of the class of all products (in our sense) of structures in K is also decidable. The same is true if we restrict ourselves to the case where the index sets are infinite and/or the filters extend the corresponding filter of cofinite sets.* \square

3. Examples

a) Let $K = \{\langle 2, \sigma \rangle\}$, where 2 is the Boolean algebra with two elements and σ is the discrete topology. Let L^{BA} be the first-order language for Boolean algebras and L_t^{BA} the topological language for Boolean algebras with topology. The L_t^{BA} -theory of K is decidable since the unique structure contained in it is finite. Thus, the class $\{\langle 2^I, \bar{\sigma}_D \mid I \text{ is a set and } D \text{ is a filter on } I\}$ has a decidable L_t^{BA} -theory. Furthermore, we shall see that it is possible to classify all its structures according to elementary equivalence.

Theorem. *For all filters D and E on I and J , respectively,*

$$\langle 2^I, \bar{\sigma}_D \rangle \equiv_{L_t^{BA}} \langle 2^J, \bar{\sigma}_E \rangle \text{ iff } \langle \mathfrak{B}(I), D \rangle \equiv_{L^{BAF}} \langle \mathfrak{B}(J), E \rangle.$$

Proof. One direction follows from the Feferman-Vaught-type theorem. The other one is obtained interpreting the L^{BAF} -theory of $\langle \mathfrak{B}(I), D \rangle$ in the L_t^{BA} -theory of $\langle 2^I, \bar{\sigma}_D \rangle$. More precisely, to each formula $\varphi(x) \in L^{BAF}$ we associate a formula $\varphi^*(x) \in L_t^{BA}$, such that, for each $I, D, a \in 2^I$:

$$\langle \mathfrak{B}(I), D \rangle \models \varphi[a] \text{ iff } \langle 2^I, \bar{\sigma}_D \rangle \models \varphi^*[a],$$

where $x = (x_1, \dots, x_n)$, $a = (a_1, \dots, a_n)$. φ^* is constructed replacing each subformula of φ of the form $t \in D$ by the L_t^{BA} -formula $(\exists X \ni t) \forall x (x \in X \rightarrow x \cap t = x)$. \square

Therefore, using the classification in elementary-equivalence types of all algebras of the form $\langle \mathfrak{B}(I), D \rangle$ made by ERSHOV [2], it is possible to classify the algebras of the form $\langle 2^I, \bar{\sigma}_D \rangle$. In particular, we can see that there are many pairs of them which are not elementary equivalent.

b) Let $(\langle A_i, \sigma_i \rangle)_{i \in I}, (\langle B_j, \tau_j \rangle)_{j \in J}$, be T_3 -spaces (regular and Hausdorff) which are not singletons. In contrast with that obtained in a), we have

Theorem. *If I, J are infinite and $D \cong Co(I)$, $E \cong Co(J)$, then*

$$\langle \prod_I A_i, \bar{\sigma}_D \rangle \equiv_{L_t} \langle \prod_J B_j, \bar{\tau}_E \rangle.$$

Proof. $\langle \prod_I A_i, \bar{\sigma}_D \rangle$ and $\langle \prod_J B_j, \bar{\tau}_E \rangle$ are T_3 -spaces without isolated points. It is known [4] that all T_3 -spaces without isolated points are L_t -equivalent. \square

c) Let us consider the non-regular T_2 topological space $\langle R, \sigma \rangle$, where the basic neighbourhoods of points different from 0 are as usual, but the basic neighbourhoods of 0 are of the form $(c, d) \setminus \{\pm 1/k \mid k \in \mathbb{N}\}$. It is possible to show that the L_t -theory of $\langle R, \sigma \rangle$ is decidable. Consequently, the theory of all topological powers $\langle R^I, \bar{\sigma}_D \rangle$ is also decidable. Not all these powers are elementary equivalent; for example, $\langle R, \sigma \rangle$ and $\langle R^I, \bar{\sigma}_D \rangle$, with $|I| > 2$, are not elementary equivalent since "there exists a unique non-regular point" is expressible in L_t . It is an open problem to classify all these powers according to elementary equivalence; in particular, for infinite index sets and filters that extend the filter of cofinite sets. In this context, it may be of interest to recall that the L_t -theory of all Hausdorff spaces is undecidable [4].

4. Properties of formulas which are preserved under topological products

A sentence $\varphi \in L_t^S$ is *preserved* (under topological products) if for all I , $(\langle \mathcal{A}_i, \sigma_i \rangle)_{i \in I}$, filter D on I : for all $i \in I$, $\langle \mathcal{A}_i, \sigma_i \rangle \models \varphi$ implies $\langle \prod_I \mathcal{A}_i, \bar{\sigma}_D \rangle \models \varphi$.

From now on, the set of symbols S will be finite.

Theorem. *The class of all preserved sentences of L_t^S is recursively enumerable.*

Proof. By the Feferman-Vaught-type theorem, we can associate effectively to each sentence $\varphi \in L_t^S$ a partitioning tuple $(\Phi; \theta_1, \dots, \theta_n)$ with $\Phi(x_0, \dots, x_{n-1}) \in L^{\text{BAF}}$ and $\theta_1, \dots, \theta_n$ L_t^S -sentences.

If $\hat{\varphi}$ is the L^{BAF} -sentence

$$\forall x_0 \dots \forall x_{n-1} \left(\bigwedge_{i < j < n} x_i \cap x_j = 0 \wedge x_0 \cup \dots \cup x_{n-1} = 1 \wedge \bigwedge_{j \in C(\varphi)} x_j = 0 \right. \\ \left. \rightarrow \Phi(x_0, \dots, x_{n-1}) \right),$$

where $C(\varphi) = \{j \mid 0 \leq j < n \text{ and } \models_{L_t^S} \varphi \rightarrow \neg \theta_j\}$, then it is easy to verify that φ is preserved iff $\hat{\varphi}$ belongs to the L^{BAF} -theory of all algebras of the form $\langle \mathfrak{P}(I), D \rangle$. As we have seen, this theory is decidable; furthermore, $\{\theta \in L_t^S \mid \models_{L_t^S} \varphi \rightarrow \neg \theta\}$ is recursively enumerable and $\hat{\varphi}$ can be obtained effectively from φ and $C(\varphi)$. From this we can conclude that the class of all preserved sentences is recursively enumerable. \square

For the usual first-order logic there is a theorem by VAUGHT [3] that establishes that a formula is preserved under direct products iff it is preserved under products of two factors. In [4] it is shown that this fact, which simplifies the study of preserved first-order formulas, does not hold for L_t and Tychonov topological products. In our case, we have the following lemma as a substitute for VAUGHT's theorem. This lemma can be proved using SKOLEM's decision method for the L^{BA} -theory of all algebras of the form $\mathfrak{P}(I)$ (see e.g. [6]).

Lemma. *To each sentence $\varphi \in L_t^S$ there can be effectively associated a natural number n_0 , such that for all $\langle \mathcal{A}, \sigma \rangle$ if $\langle \mathcal{A}, \sigma \rangle^{n_0} \models \varphi$, then for all $m \geq n_0$ $\langle \mathcal{A}, \sigma \rangle^m \models \varphi$. ($\langle \mathcal{A}, \sigma \rangle^m$ is the usual topological power). \square*

From this lemma and a series of technical results with classical analogues, it follows, as in [5] for first-order logic and reduced products, the following theorem:

Theorem. *Each sentence of L_t^S is logically equivalent to a Boolean combination of sentences which are preserved under topological products with respect to non-trivial filters. \square*

In what follows, we assume that the filters extend the corresponding filter of cofinite sets. The next two theorems give us a large syntactical class of formulas of L_t which are preserved. The first one can be proved by induction.

Theorem.

- a) *Atomic formulas and negations of atomic formulas are preserved.*
- b) *If φ is preserved, then $\exists x\varphi$ and $\forall x\varphi$ are preserved.*
- c) *If φ and ψ are preserved, then $\varphi \wedge \psi$ is preserved.*
- d) *If φ is preserved and positive in X , then $(\forall X \ni t)\varphi$ is preserved.*

Theorem. *If $\varphi \in L_t^S$ is positive, without second-order quantifiers and Q is a block of second-order existential quantifiers, then the following formulas of L_t^S are preserved:*

- a) *$\varphi \rightarrow \theta$ if φ is preserved.*
- b) *$Q(\varphi \rightarrow \theta)$ if θ is preserved and positive.*
- c) *$Q\forall y(\varphi \rightarrow \theta)$ if θ is a second-order atomic formula, i.e. of the form $t \in X$.*

Proof.

a) It is possible to show that the formulas of L_t^S which are preserved under continuous surjective homomorphisms are those that are equivalent to positive formulas without second-order existential quantifiers. On the other hand, by the hypothesis on the filters, the projection maps are continuous. The combination of these facts gives the proof.

b) As $Q(\varphi \rightarrow \theta) \in L_t^S$, the second-order variables in Q do not appear in θ ; hence $Q(\varphi \rightarrow \theta) \vdash Q'\varphi \rightarrow \theta$, where Q' is a block of second-order universal quantifiers. Therefore we can apply a).

c) For simplicity, we restrict ourselves to the case $Q = (\exists Y \ni t)$, $y = y$, that is, we have to prove that

$$\psi := (\exists Y \ni t(x, \dots)) \forall y(\varphi(y, x, \dots, Y, X, \dots) \rightarrow \theta(y, x, \dots, X))$$

is preserved.

Let us suppose that for all $i \in I$

$$\langle \mathcal{A}_i, \sigma_i \rangle \models \psi(x, \dots, X, \dots) [b_i, \dots, U_i, \dots]$$

(here $b, \dots \in \prod_I A_i$, $\prod_I U_i, \dots \in \sigma_D$). Then, for each $i \in I$, let $Y \in \sigma_i$ be such that $t^{\mathcal{A}_i}[b_i, \dots] \in Y_i$ and

$$(1) \quad \langle \mathcal{A}_i, \sigma_i \rangle \models \forall y(\varphi(y, x, \dots, Y, X, \dots) \rightarrow \theta(y, x, \dots, X)) [b_i, \dots, Y_i, U_i, \dots].$$

If, for all $i \in I$, one has $\langle \mathfrak{A}_i, \sigma_i \rangle \models \forall y \theta(y, x, \dots, X) [b_i, \dots, U_i]$, then, by the preceding theorem, $\langle \prod_I \mathfrak{A}_i, \bar{\sigma}_D \rangle \models \forall y \theta(y, x, \dots, X) [b, \dots, \prod_I U_i]$. Taking $Y = \prod_I A_i$, it follows that $\langle \prod_I \mathfrak{A}_i, \bar{\sigma}_D \rangle \models \psi[b, \dots, \prod_I U_i, \dots]$.

Otherwise, there exist sets $I_0 \subseteq I_1 \in D$ such that

$$i_0 \in I_0 \text{ iff } (\mathfrak{A}_{i_0}, \sigma_{i_0}) \models \exists y \neg \theta[b_{i_0}, \dots, U_{i_0}].$$

Then, for all $i_0 \in I_0$ and all $a_{i_0} \in A_{i_0}$ such that $\langle \mathfrak{A}_{i_0}, \sigma_{i_0} \rangle \models \neg \theta[a_{i_0}, b_{i_0}, \dots, U_{i_0}]$, one has by (1) that $\langle \mathfrak{A}_{i_0}, \sigma_{i_0} \rangle \models \neg \varphi(y, x, \dots, Y, X) [a_{i_0}, b_{i_0}, \dots, Y_{i_0}, U_{i_0}, \dots]$.

As $\neg \varphi$ is negative without second-order universal quantifiers, $\neg \varphi$ is preserved under preimages of continuous surjective homomorphisms, hence

$$\langle \prod_I \mathfrak{A}_i, \bar{\sigma}_D \rangle \models \neg \varphi(y, x, \dots, Y, X, \dots) [\tilde{a}_{i_0}, \tilde{b}_{i_0}, \dots, p_{i_0}^{-1}(Y_{i_0}), p_{i_0}^{-1}(U_{i_0}), \dots],$$

where $\tilde{a}_{i_0}, \tilde{b}_{i_0}, \dots \in \prod_I A_i$ are preimages of a_{i_0}, b_{i_0}, \dots under $p_{i_0}^{-1}$ (in particular, we can take $\tilde{b}_{i_0}, \dots = b, \dots$). In consequence, letting $V := \bigcap_{i \in I_1} p_i^{-1}(Y_i) \subseteq \bigcap_{i \in I_0} p_i^{-1}(Y_i) \subseteq p_{i_0}^{-1}(Y_{i_0})$, we obtain

$$(2) \quad \langle \prod_I \mathfrak{A}_i, \bar{\sigma}_D \rangle \models \neg \varphi[\tilde{a}_{i_0}, b, \dots, V, \prod_I U_i, \dots].$$

Notice that $V \in \sigma_D$ and therefore we can replace $p_{i_0}^{-1}(Y_{i_0})$ by V because $\neg \varphi$ is negative.

Now let us suppose that $\langle \prod_I \mathfrak{A}_i, \bar{\sigma}_D \rangle \not\models \psi[b, \dots, \prod_I U_i, \dots]$; then

$$\begin{aligned} \langle \prod_I \mathfrak{A}_i, \bar{\sigma}_D \rangle \models (\forall Y \ni t(x, \dots)) \exists y (\varphi(y, x, \dots, Y, X, \dots) \wedge \\ \wedge \neg \theta(y, x, \dots, X)) [b, \dots, \prod_I U_i, \dots]. \end{aligned}$$

In particular, taking the above defined V as assignment for Y , we obtain an $a \in \prod_I A_i$ such that

$$(3) \quad \begin{aligned} \langle \prod_I \mathfrak{A}_i, \bar{\sigma}_D \rangle \models \varphi(y, x, \dots, Y, X, \dots) [a, b, \dots, V, \prod_I U_i, \dots] \text{ and} \\ \langle \prod_I \mathfrak{A}_i, \bar{\sigma}_D \rangle \models \neg \theta(y, x, \dots, X) [a, b, \dots, \prod_I U_i]. \end{aligned}$$

From this last fact there follows the existence of $i_0 \in I_0$ such that $\langle \mathfrak{A}_{i_0}, \sigma_{i_0} \rangle \models \neg \theta(y, x, \dots, X) [a_{i_0}, b_{i_0}, \dots, U_{i_0}]$. As we now are in the situation that led us to (2), we conclude that

$$\langle \prod_I \mathfrak{A}_i, \bar{\sigma}_D \rangle \models \neg \varphi(y, x, \dots, Y, X, \dots) [\tilde{a}_{i_0}, b, \dots, V, \prod_I U_i, \dots].$$

In particular, with $\tilde{a}_{i_0} = a$, it follows

$$\langle \prod_I \mathfrak{A}_i, \bar{\sigma}_D \rangle \models \neg \varphi(y, x, \dots, Y, X, \dots) [a, b, \dots, V, \prod_I U_i, \dots],$$

which contradicts (3). \square

Applying systematically the two preceding theorems it is possible to verify syntactically that the following topological properties expressible in L_t are preserved:

$$\forall x \forall y ((\forall X \ni x) y \in X \wedge (\forall Y \ni y) x \in Y) \rightarrow x = y \quad (``T_0"),$$

$$\forall x \forall y ((\forall X \ni x) y \in X \rightarrow x = y) \quad (``T_1"),$$

$$\forall x \forall y (\exists X \ni x) (\exists Y \ni y) (\exists z (z \in X \wedge z \in Y) \rightarrow x = y) \quad (``Hausdorff"),$$

$$\begin{aligned}
&\forall x(\forall X \ni x) (\exists Y \ni x) \forall y((\forall W \ni y) \exists z(z \in W \wedge z \in Y) \rightarrow y \in X) && \text{("regular")}, \\
&\forall x(\forall X \ni f(x)) (\exists Y \ni x) \forall z(z \in Y \rightarrow f(z) \in X) && \text{("f is continuous")}, \\
&\forall x((\forall X \ni x) \exists y(Py \wedge y \in X) \rightarrow Px) && \text{("P is closed")}.
\end{aligned}$$

References

- [1] BERTOSSI, L. E., El lenguaje formal L_t y productos topológicos. Tesis de Doctorado en Matemática, Pontificia Universidad Católica de Chile, 1987.
- [2] ERSHOV, Y. L., Theorie der Numerierungen III. This Zeitschrift **23** (1977), 289–371.
- [3] FEFERMAN, S., and R. L. VAUGHT, The first order properties of products of algebraic systems. Fund. Math. **47** (1959), 57–103.
- [4] FLUM, J., and M. ZIEGLER, Topological Model Theory. Springer Lecture Notes in Mathematics **769** (1980).
- [5] GALVIN, F., Horn sentences. Annals Math. Logic **1** (1970), 389–422.
- [6] GRAETZER, G., Universal Algebra. Springer-Verlag, Berlin-Heidelberg-New York 1979.
- [7] ZIEGLER, M., Topological model theory. In: Model Theoretic Logics (J. BARWISE and S. FEFERMAN, eds.), Springer-Verlag, Berlin-Heidelberg-New York 1985, pp. 557–577.

L. E. Bertossi

(Eingegangen am 14. Oktober 1988)

Pontificia Universidad Católica de Chile
 Facultad de Matemática
 Casilla 6177
 Santiago 22
 Chile