

School of Computer Science

Carleton University

Name:
Course: COMP 3804

Problem Set: 1
Week: 1

Problems

Consult references [1, 2].

Problem 1. Define O, Ω, Θ in the context of this course.

Problem 2. Applying the definitions, show the following:

1. $1 + 2 + 3 + \dots + n = \frac{n(n+1)}{2}$.
2. $1 + 2 + 3 + \dots + n \in O(n^2)$.
3. $1 + 2 + 3 + \dots + n \in \Omega(n^2)$.
4. $1 + 2 + 3 + \dots + n \in \Theta(n^2)$.

Problem 3. Show that $\sum_{k=1}^n k^2 \in \Theta(n^3)$

Problem 4. Show that $\frac{1}{3}n^2 - 3n \in O(n^2)$

Problem 5. Let $p(n) = \sum_{i=0}^d a_i n^i$ be a polynomial of degree d and assume that $a_d > 0$. Show that $p(n) \in O(n^k)$, where $k \geq d$ is a constant. What are c and n_0 if we use the following definition of O -notation: $O(g(n)) = \{f(n) : \exists c, n_0 > 0 \text{ such that } f(n) \leq cg(n), \forall n \geq n_0\}$.

Problem 6. Let $p(n) = \sum_{i=0}^d a_i n^i$ be a polynomial of degree d and assume that $a_d > 0$. For $k \geq d$, show that $\lim_{n \rightarrow \infty} \frac{\sum_{i=0}^d a_i n^i}{n^k} = a_d n^{d-k} \geq a_d > 0$ and conclude that $p(n) \in \Omega(n^k)$.

Problem 7. Let $p(n) = \sum_{i=0}^d a_i n^i$ be a polynomial of degree d and assume that $a_d > 0$. Show that $p(n) \in \Theta(n^d)$. (Hint: Using limits show that $\lim_{n \rightarrow \infty} \frac{p(n)}{n^d} = a_d$ and $0 < a_d < \infty$. Thus, $p(n) \in \Theta(n^d)$. Recall that $f(n) \in \Theta(g(n))$, if $0 < \lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} < \infty$.)

Problem 8. Show that $6n \log n + \sqrt{n} \log^2 n = \Theta(n \log n)$.

Problem 9. Fibonacci numbers are defined recursively as follows:

$F_0 = 0, F_1 = 1$, and $F_n = F_{n-1} + F_{n-2}$ for any integer $n > 1$.

Using induction show that $F_n \geq 2^{\frac{n}{2}}$ for any $n \geq 6$.

References

- [1] S. DasGupta, C. Papadimitriou, V. Vazirani. *Introduction to Algorithms*. McGraw Hill.
- [2] A. Maheshwari. *Notes on Algorithm Design*, Chapter 1, <https://people.scs.carleton.ca/~maheshwa/Notes/DAA/notes.pdf>