TARANET: Traffic-Analysis Resistant Anonymity at the Network Layer

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EuroS&P 2018



Our vision:

Our vision:

Anonymous communication

Anonymous communication Non-discrimination

Our vision:

Our vision:

Anonymous communication

Non-discrimination

Prevent industrial espionage

















































































Performance









Performance



Hsiao et al. LAP: Lightweight Anonymity and Privacy. In IEEE S&P, 2012 ullet







Performance



- \bullet
- Sankey, Wright. *Dovetail: Stronger anonymity in next-generation internet routing*. In PETS, 2014







Performance



- \bullet
- Chen, Perrig. PHI: Path-Hidden Lightweight Anonymity Protocol at Network Layer. In PETS, 2017 \bullet







Performance



- \bullet
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Traffic analysis Severe threat to anonymous communication



anonymity



























Starting point

Layered encryption







Starting point

- Layered encryption
- Per-hop authentication







Starting point

- Layered encryption
- Per-hop authentication
- Fixed packet length






































































Packet counting











THzürich

- Inter-packet timing















































































TARANET Resisting traffic analysis attacks























• Fixed rate 1/B









- Fixed rate 1/B
- Chaff packets









- Fixed rate 1/B
- Chaff packets
- Fixed duration T









- Fixed rate 1/B
- Chaff packets
- Fixed duration T









- Fixed rate 1/B
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- Fixed rate 1/B
- Chaff packets
- Fixed duration T







- Fixed rate 1/B
- Chaff packets
- Fixed duration T
- Same for everyone









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- Fixed rate 1/B
- Chaff packets
- Fixed duration T
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Flowlet









- Fixed rate 1/B
- Chaff packets
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- Same for everyone



Flowlet







- Fixed rate 1/B
- Chaff packets
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Flowlet







- Fixed rate 1/B
- Chaff packets
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- Same for everyone


































Resisting passive traffic analysis *re-enforce* flowlet schedule







Resisting passive traffic analysis *re-enforce* flowlet schedule







Resisting passive traffic analysis *re-enforce* flowlet schedule U







Resisting passive traffic analysis *re-enforce* flowlet schedule























Resisting active traffic analysis **Packet dropping** U









Resisting active traffic analysis **Packet dropping** U



















Duplicate packets









- Duplicate packets
- Create new packets











- Duplicate packets
- Create new packets
- Sender send more?





































































When is packet splitting done?









When is packet splitting done?

Sender includes splittable packets









When is packet splitting done?

- Sender includes splittable packets
- ... for each AS on the path









When is packet splitting done?

- Sender includes splittable packets
- ... for each AS on the path
- ... at random intervals







How does splitting work concretely?







How does splitting work concretely?







| R | MAC | R | MAC | MAC | |
|---|-----|---|-----|-----|--|
| | | | | | |





Payload











Payload













Payload











Payload







| MAC |
|-----|
|-----|





Payload







| | R | MAC | R | MAC | |
|--|---|-----|---|-----|--|
|--|---|-----|---|-----|--|





| | R | MAC | R | MAC | |
|--|---|-----|---|-----|--|
|--|---|-----|---|-----|--|









































Pseudorandom Payload

Pseudorandom Payload













Pseudorandom Payload

Pseudorandom Payload



TARANET PerformanceEvaluation setup, Throughput, Latency





Evaluation setup





- Prototype implementation
 - Data-Plane Development Kit (DPDK)






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 - Data-Plane Development Kit (DPDK)
- Software router
 - 12x 10 GbE NICs
 - Intel Xeon 2.7 GHz (2x 8 cores)







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 - Single 10 GbE interface







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 - Single 10 GbE interface
 - Single processing core



Throughput





| 768 | 1024 | 1280 |
|-----------------|------|------|
| ad Size (bytes) | | |





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Throughput

| 768 | 1024 | 1280 |
|-----------------|------|------|
| ad Size (bytes) | | |











| 768 | 1024 | 1280 |
|-----------------|------|------|
| ad Size (bytes) | | |





Path length: 7 nodes





| 768 | 1024 | 1280 |
|-----------------|------|------|
| ad Size (bytes) | | |





Path length: 7 nodes









Path length: 7 nodes









Path length: 7 nodes











Latency

| 768 | 1024 | 1280 |
|-----------------|------|------|
| ad Size (bytes) | | |







Latency



Summary Highlights, Limitations



- TARANET highlights:
 - Protection against passive traffic analysis with *flowlets*
 - Protection against active traffic analysis with packet splitting





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 - Good performance 3 Gbps on single core, acceptable latency





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- Limitations: lacksquare





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- Limitations:
 - Chaff traffic creates a non-negligible bandwidth overhead





- TARANET highlights:
 - Protection against passive traffic analysis with *flowlets*
 - Protection against active traffic analysis with packet splitting
 - Good performance 3 Gbps on single core, acceptable latency
- Limitations:
 - Chaff traffic creates a non-negligible bandwidth overhead
 - Third-party anonymity





In the paper

TARANET: Traffic-Analysis Resistant Anonymity at the Network Layer

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Abstract—Modern low-latency anonymity systems, no matter whether constructed as an overlay or implemented at the network layer, offer limited security guarantees against traffic analysis. On the other hand, high-latency anonymity systems offer strong security guarantees at the cost of computational overhead and long delays, which are excessive for interactive applications. We propose TARANET, an anonymity system that implements protection against traffic analysis at the network layer, and limits the incurred latency and overhead. In TARANET's setup phase, traffic analysis is thwarted by mixing. In the data transmission phase, end hosts and ASes coordinate to shape traffic into constant-rate transmission using packet splitting. Our prototype implementation shows that TARANET can forward anonymous traffic at over 50 Gbps using commodity hardware.

1. Introduction

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Users are increasingly aware of their lack of privacy and

in forwarding anonymous traffic. Intermediate anonymity supporting network nodes (or nodes for short) first cooperate with senders to establish anonymous sessions or circuits, and then process and forward traffic from those senders to receivers. While these systems achieve high throughput and low latency, the security guarantees of these systems are no stronger than Tor's. Moreover, LAP and Dovetail leak the position of intermediate nodes on the path and the total path length, which reduces the anonymity set size, facilitating deanonymization [21].

The problem space appears to have an unavoidable tradeoff: strong anonymity appears achievable only through *drastically higher overhead* [27]. In this paper, we aim to push the boundaries of this anonymity/performance tradeoff by combining the speed of network-layer anonymity systems with strong defenses.

To improve the anonymity guarantees, traffic analysis attacks need to be prevented, or made significantly harder for the adversary to perform. The common method to achieve

- Flowlet setup (asymm. crypto)
- Link padding (security in depth)
- Anonymity set size analysis
- Security analysis

- Chaff/setup packet trade-off
- **Deployment incentives**





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Thank you!