

Probabilistic Method

- MAX CUT
- RAMSEY NUMBERS $R(k, k)$
- HAMILTONIAN PATHS IN TOURNAMENTS
- Sperner's Thm

Probabilistic Method

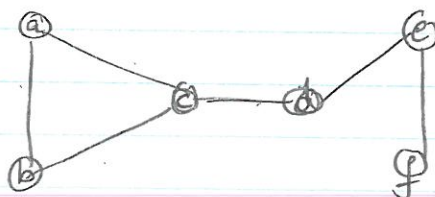
Let $G = (V, E)$ be a simple graph.

↑ ↑
vertices edges.

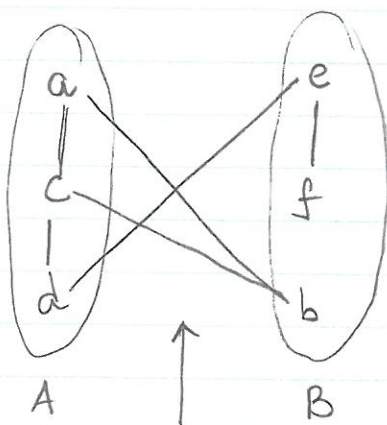
Let $V = A \cup B$ and $A \cap B = \emptyset$.

We say an edge $e \in E$ is between A and B if one of its endpoints is in A and other is in B .

Consider the following graph



Define A & B such that



In between edges

Notice that in this case at least half of the edges of G are "between" edges.

Claim: Let $G=(V,E)$ be a simple graph with m edges. The vertex set V can be partitioned into two subsets A and B such that the number of edges between A and B is at least $m/2$.

Proof: Take a fair coin and toss it $|V|$ times, once for each vertex. If the outcome is Heads, add the vertex $v \in V$ to A , otherwise add it to B .

Define r.v. $X = \#$ edges between A & B .

Let $E = (e_1, e_2, \dots, e_m)$ be the edges of G .

Define m 0-1 r.v. $X_1, X_2, \dots, X_m$ as follows

$$X_i = \begin{cases} 1 & \text{if } e_i \text{ is an edge between } A \text{ \& } B \\ 0 & \text{otherwise} \end{cases}$$

$$X = \sum_{i=1}^m X_i$$

$$E(X) = E\left(\sum_{i=1}^m X_i\right) = \sum_{i=1}^m E(X_i) = \sum_{i=1}^m \Pr(X_i=1)$$

$G(n, p)$ properties:

1. $np < 1$ No connected components of size larger than $O(\log n)$
2. $np = 1$ largest conn comp of size $n^{2/3}$
3. $np = c > 1$ unique giant component with a fraction of vertices - no other components have more than $O(\log n)$ vertices
4. $p < \frac{(1-\epsilon) \ln n}{n}$ $G(n, p)$ will have isolated vertices
5. $p > \frac{(1+\epsilon) \ln n}{n}$ $G(n, p)$ will be connected.

Let the endpoints of $e_i = (u, v)$ where $u, v \in V$.

Case 1:	$u \in A, v \in A$	$X_i = 0$
Case 2:	$u \in A, v \in B$	$X_i = 1$
Case 3:	$u \in B, v \in A$	$X_i = 1$
Case 4:	$u \in B, v \in B$	$X_i = 0$

Since $X_i = 1$ in two out of four cases, and each case is equally likely to occur, thus we have

$$E(X_i) = \Pr(X_i = 1) = 1/2.$$

$$\Rightarrow E(X) = \Pr(X_1 = 1) + \Pr(X_2 = 1) + \dots + \Pr(X_m = 1) \\ = m/2. \quad \text{--- } (***)$$

Suppose that the claim in theorem is false.

$\Rightarrow$ No matter how we partition the vertices

in sets $A \neq B$, $A \cap B = \emptyset$ and $A \cup B = V$,

the # in between edges $< \frac{m}{2}$.

$\Rightarrow E(X) < \frac{m}{2}$ and that contradicts $(***)$

Comment: 2 approx to MAX-CUT

Example 2:

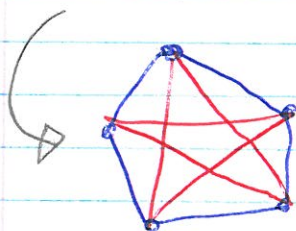
RAMSEY THEORY.

Recall that $R(k, l)$ for integers $k \geq 0, l \geq 0$ is defined to be

$$R(k, l) = \min \{ n \mid \text{any graph on } n \text{ vertices whose edges are colored red or blue always contains either a red clique of size } k \text{ or a blue clique of size } l \}.$$

In the very first class we saw that

$$R(3, 3) > 5 \text{ and } R(3, 3) = 6$$



No red or blue triangle

Claim: For any $k \geq 3$, $R(k, k) > \left\lfloor 2^{k/2} \right\rfloor$.

Proof: Consider a complete graph on n vertices, and color each edge red or blue with $p_r = 1/2$, independent of any other edge.

(e.g. flip a fair coin for each edge

"heads" $\rightarrow$ color it red

"tails" $\rightarrow$ color it blue)

Let $G=(V, E)$ be the complete bichromatic graph.

Take a set of k vertices, $S \subseteq V$.

$$\Pr(\text{all edges in the graph induced by } S \text{ are red}) = \left(\frac{1}{2}\right)^{\binom{k}{2}}$$

$$\Pr(\text{all edges in the graph induced by } S \text{ are blue}) = \left(\frac{1}{2}\right)^{\binom{k}{2}}$$

In all there are $\binom{n}{k}$ possible subsets of k vertices out of n vertices.

Thus

$$\Pr(G \text{ contains a red clique or a blue clique of size } k) \leq 2 \cdot \left(\frac{1}{2}\right)^{\binom{k}{2}} \binom{n}{k}$$

This follows from union bound.

What value of n ensures that

$$\Pr(\text{)} \leq 1 \text{ ?}$$

Note that $\binom{n}{k} \leq n^k$

$$\text{If } 2 \binom{n}{k} \leq n^k < 2^{k(k-1)/2} \text{ then } \Pr(\text{)} < 1.$$

$$\text{i.e. if } n \leq 2^{k/2-1} \text{ then}$$

$$2 \cdot \left(2^{k/2-1}\right)^k = 2 \cdot 2^{k(k/2-1)} = 2^{k(k/2)} = 2^{k^2/2}$$

Consider the quantity

$$2 \cdot \left(\frac{1}{2}\right)^{\binom{k}{2}} \binom{n}{k}$$

$$= \frac{2 \binom{n}{k}}{2^{\binom{k}{2}}} \quad \text{--- (1)}$$

We are interested to know what value of n will ensure that (1) is < 1 .

$\Rightarrow \Pr(G \text{ does not contain a red clique of size } k \text{ and } G \text{ does not contain a blue clique of size } k) > 0.$

$\rightarrow$ Set $n = \lfloor 2^{k/2} \rfloor \leq 2^{k/2}$

Then (1) becomes

$$\frac{2 \binom{n}{k}}{2^{\binom{k}{2}}} = \frac{n(n-1)(n-2) \cdots (n-k+1)}{k!} \cdot \frac{2}{2^{(k^2-k)/2}}$$

$$\leq \frac{n^k}{k!} \cdot \frac{2^{1+k/2}}{2^{k^2/2}}$$

$$\leq \frac{(2^{k/2})^k}{k!} \cdot \frac{2^{1+k/2}}{2^{k^2/2}} = \frac{2^{1+k/2}}{k!}$$

Example 3: Hamiltonian Paths.

A Hamiltonian Path in a directed graph G is a directed path that passes through all vertices of G .

Claim: Each edge of a complete graph on n -vertices can be directed in such a way that G has at least $\frac{n!}{2^{n-1}}$ directed Hamiltonian paths.

Pf: Consider a complete graph $G = (V, E)$ on n -vertices.

For each edge $e = (u, v) \in E$, $u \in V, v \in V$, toss a fair coin (independently of other edges).

If the coin comes up heads orient e from u to v .
tails orient e from v to u .

Consider a permutation of n vertices — and let that permutation be $\alpha = (\alpha(1), \alpha(2), \dots, \alpha(n))$.

If in G , each edge $\alpha(i)\alpha(i+1)$ for $1 \leq i \leq n-1$, appears in that orientation, then

$(\alpha(1), \alpha(2), \dots, \alpha(n))$ is a Hamiltonian path.

Let $X_\alpha = \begin{cases} 1 & \alpha \text{ is a Hamiltonian path.} \\ 0 & \text{otherwise.} \end{cases}$

$$E[X_\omega] = \Pr \left(\text{Each edge } \sigma(i) \rightarrow \sigma(i+1) \text{ is oriented} \right. \\ \left. \text{in this way for } \forall 1 \leq i \leq n-1 \right)$$

$$= \frac{1}{2^{n-1}} \quad \text{as each edge orientation is} \\ \text{chosen independently.}$$

Thus total # of Hamiltonian paths X equals to

$$E[X] = E \left[\sum_{\omega \in \text{permutation of } n} X_\omega \right] = \sum_{\omega} E[X_\omega]$$

$$= \frac{n!}{2^{n-1}} \quad \blacksquare$$

If expected # of directed Hamiltonian paths in a "randomly-oriented graph" is

$\frac{n!}{2^{n-1}}$, then there is always a way to

orient the graph such that # directed Hamiltonian paths $\geq \frac{n!}{2^{n-1}}$.

Sperner's Theorem

Let $n \geq 1$ be an integer and let $S = \{1, 2, 3, \dots, n\}$.

Let $S_1, S_2, \dots, S_m$ be a sequence of m subsets of S such that

$$\forall i, j, i \neq j, S_i \not\subseteq S_j \text{ and } S_j \not\subseteq S_i.$$

Then $m \leq \binom{n}{\lfloor n/2 \rfloor}. \quad \text{--- (1)}$

Proof: Assume none of the subsets $S_1, S_2, \dots, S_m$

is empty, otherwise $m=1$ and (1) holds trivially.

Choose a random permutation of elements of S .

e.g. for $n=4$, let a permutation be $\{3, 1, 4, 2\}$.

Define sets $A_1 = \{3\}$
 $A_2 = \{1, 3\}$
 $A_3 = \{1, 3, 4\}$
 $A_4 = \{1, 2, 3, 4\}$

} Random subsets of S as they are based on a random permutation

If any of $\{S_1, S_2, \dots, S_m\}$ equals any of $\{A_1, A_2, \dots, A_n\}$ we say ^{"n"} some S_i occurs in the ["]sequence $A_1, A_2, \dots, A_n$.

Clearly, at most one S_i can occur in

$A_1, A_2, \dots, A_n$ as $A_1 \subseteq A_2 \subseteq A_3 \subseteq \dots \subseteq A_n$.

Define r.v. $X = \begin{cases} 1 & \text{if any of } S_i \text{ occurs in } A_1, A_2, \dots, A_n \\ 0 & \text{otherwise} \end{cases}$

$$E(X) \leq 1.$$

For each $1 \leq i \leq m$, define

$$X_i = \begin{cases} 1 & \text{if } S_i \text{ occurs in sequence } A_1, A_2, \dots, A_n \\ 0 & \text{otherwise.} \end{cases}$$

$$\text{Let } k = |S_i|$$

Note that $X_i = 1$ iff $S_i = A_i$.

$$\text{Let } A_k = \{a_1, a_2, \dots, a_k\} \in S.$$

Hence, $X_i = 1$ iff the first k values in the random permutation form a permutation of the elements of S_i

$$\begin{array}{c} \boxed{a_1, a_2, \dots, a_k \mid a_{k+1} \dots a_n} \\ \underbrace{\hspace{1.5cm}} \\ \text{permutation} \\ \text{of } S_i \end{array}$$

In all there are $k!(n-k)!$ total # of permutations of n -elements, where first k elements correspond to the set S_i .

$$\text{Thus, } E(X_i) = \Pr(X_i = 1)$$

$$= \frac{k!(n-k)!}{n!} = \frac{1}{\binom{n}{k}} = \frac{1}{\binom{n}{|S_i|}}$$

Now

$$X = \sum_{i=1}^m X_i$$

$$E(X) = E\left(\sum_{i=1}^m X_i\right)$$

$$= \sum_{i=1}^m E(X_i)$$

$$= \sum_{i=1}^m \frac{1}{\binom{n}{|S_i|}}$$

Also, we know $E(X) \leq 1$.

Thus

$$\sum_{i=1}^m \frac{1}{\binom{n}{|S_i|}} \leq 1$$

Note that $\binom{n}{|S_i|} \leq \binom{n}{\lfloor \frac{n}{2} \rfloor}$ as largest value in a binomial is in the middle (see Pascal's Δ)

Hence $1 \geq \sum_{i=1}^m \frac{1}{\binom{n}{|S_i|}} \geq \sum_{i=1}^m \frac{1}{\binom{n}{\lfloor \frac{n}{2} \rfloor}} = \frac{m}{\binom{n}{\lfloor \frac{n}{2} \rfloor}}$

or $\binom{n}{\lfloor \frac{n}{2} \rfloor} \geq m$