## Balls \& Bins

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## Outline

Model

Collisions

Size of Bins

Model

## Balls \& Bins

## Model

We have $m$ Balls and $n$ Bins. We throw each ball in a bin uniformly at random.

What is the probability of following events:

1. Balls $i$ and $j$ are in the same bin.
2. Bin $\# i$ receives (a) 0 balls, (b) $k$ balls, and (c) $\geq k$ balls.
3. All bins have $\leq \frac{c \ln n}{\ln \ln n}$ balls.

Applications: Birthday Paradox, Load Balancing, Perfect Hashing

## Collisions

## Probability of Balls $i$ and $j$ in the same bin

Number of Balls $=m$
Number of Bins $=n$.
$\operatorname{Pr}[$ Balls $i$ and $j$ in same bin $]=\frac{1}{n}$.

## Expected number of collisions

Number of Balls $=m$
Number of Bins $=n$.
Define r.v. $X_{i j}(1 \leq i \leq m-1, i+1 \leq j \leq m)$ as follows:

$$
X_{i j}= \begin{cases}1 & \text { if balls } i \text { and } j \text { are in same bin } \\ 0 & \text { Otherwise }\end{cases}
$$

$\operatorname{Pr}\left(X_{i j}=1\right)=\frac{1}{n}$ and $E\left[X_{i j}\right]=\frac{1}{n}$
Define $X=\sum_{i, j} X_{i j}=$ Total \# of collisions
$E[X]=E\left[\sum_{i, j} X_{i j}\right]$
By Linearity of Expectation: $E\left[\sum_{i, j} X_{i j}\right]=\sum_{i, j} E\left[X_{i j}\right]$
Thus $E[X]=\frac{1}{n}\binom{m}{2}$

## Birthday Paradox

Number of Balls $=m=$ Number of Students
Number of Bins $=n=$ Number of days in a Year.
For two students to have same Birthday:
What value of $m$ will result in $E[X]=\frac{1}{n}\binom{m}{2} \geq 1$
For $m \geq 28, E[X]=\frac{1}{365} \frac{m(m-1)}{2} \geq 1$

## Birthday Paradox Contd.

What is minimum value of $m$ so that the probability that two students share the same birthday is $\geq \frac{1}{2}$ ?

Probability that all $m$ students have distinct birthday's is given by: $\left(1-\frac{1}{n}\right)\left(1-\frac{2}{n}\right)\left(1-\frac{3}{n}\right) \ldots\left(1-\frac{m-1}{n}\right)$

Let us use the inequality $1-x \leq e^{-x}$.
We want:

$$
\begin{aligned}
e^{-\frac{1}{n}} e^{-\frac{2}{n}} e^{-\frac{3}{n}} \ldots e^{-\frac{m-1}{n}} & \leq \frac{1}{2} \\
e^{-\frac{m(m-1)}{2 n}} & \leq \frac{1}{2}
\end{aligned}
$$

Now using $n=365$, we have $e^{-\frac{m(m-1)}{2 * 365}} \leq \frac{1}{2}$ or $m \geq 23$.

## Size of Bins

## Number of Balls in Bin $i$

Number of Balls $=m$; Number of Bins $=n$.

## Problem I

What is the probability that Bin $i$ receives no balls?

$$
\begin{aligned}
& \qquad\left(1-\frac{1}{n}\right)^{m} \leq e^{-\frac{m}{n}} \\
& \text { If } n=m,\left(1-\frac{1}{n}\right)^{n} \leq e^{-1}=0.37 .
\end{aligned}
$$

## Problem II

What is the probability that Bin $i$ receives exactly $k$ balls?

$$
\binom{m}{k}\left(\frac{1}{n}\right)^{k}\left(1-\frac{1}{n}\right)^{m-k}
$$

## Number of Balls in Bin $i$ contd.

Number of Balls $=m$; Number of Bins $=n$.

## Problem III

What is the probability that Bin $i$ receives $\geq k$ balls?

$$
\leq\binom{ m}{k}\left(\frac{1}{n}\right)^{k}
$$

If $n=m$ and using Stirling's approximation $\left.\binom{n}{k} \leq\left(\frac{e n}{k}\right)^{k}\right)$, we have $\binom{n}{k}\left(\frac{1}{n}\right)^{k} \leq\left(\frac{e}{k}\right)^{k}$

## Expected Number of Balls in a Bin

Number of Balls $=m$; Number of Bins $=n$.

## Problem IV

## What is Expected \# of Balls in a Bin?

$$
=\frac{m}{n}
$$

Define a r.v. $B_{i j}$ such that

$$
B_{i j}= \begin{cases}1 & \text { if balls } i \text { is in } \operatorname{Bin} j \\ 0 & \text { Otherwise }\end{cases}
$$

Result follows from $\operatorname{Pr}\left(B_{i j}=1\right)=\frac{1}{n}$ and $E\left[B_{i j}\right]=\frac{1}{n}$.

## Expected Number Empty Bins

Number of Balls $=m$; Number of Bins $=n$.

## Problem V

## What is Expected \# of Empty Bins?

Define a r.v. $X_{i}$ such that

$$
X_{i}= \begin{cases}1 & \text { if Bin } i \text { is empty } \\ 0 & \text { Otherwise }\end{cases}
$$

From Problem I, $\operatorname{Pr}\left(X_{i}=1\right) \leq e^{-\frac{m}{n}}$ and $E\left[X_{i}\right] \leq e^{-\frac{m}{n}}$
Thus, $E[\#$ of Empty Bins $]=\sum_{i=1}^{n} E\left[X_{i}\right] \leq n e^{-\frac{m}{n}}$
When $n=m, E[\#$ of Empty Bins $] \leq \frac{n}{e}$

## Max \# Balls in Bins

Number of Balls $=$ Number of Bins $=n$.

## Max \# of Balls in Bins

With probability $\geq 1-\frac{1}{n}$ all bins receive fewer than $3 \frac{\ln n}{\ln \ln n}$ balls.
$\operatorname{Pr}(\operatorname{Bin} i$ has more that $k$ balls $) \leq\left(\frac{e}{k}\right)^{k}$
Substitute $k=3 \frac{\ln n}{\ln \ln n}$ and show that
$\operatorname{Pr}(\operatorname{Bin} i$ has more that $k$ balls $) \leq \frac{1}{n^{2}}$
Thus by Union Bound,
$\operatorname{Pr}($ Any bin has more that $k$ balls $) \leq \frac{1}{n}$

## Max \# Balls in Bins Contd.

## Claim

For $k=3 \frac{\ln n}{\ln \ln n},\left(\frac{e}{k}\right)^{k} \leq \frac{1}{n^{2}}$

## Proof.

$$
\begin{aligned}
\left(\frac{e}{k}\right)^{k} & =\left[e^{\ln \ln n} 3 \ln \right]^{\frac{3 \ln n}{\ln \ln n}} \\
& =\left[e^{1} e^{\left.\ln \frac{\ln \ln n}{3 \ln n}\right]^{\frac{3 \ln n}{\ln \ln n}}}\right. \\
& =e^{\frac{3 \ln n}{\ln \ln n}[1+\ln \ln \ln n-\ln (3 \ln n)]} \\
& =e^{\frac{3 \ln n}{\ln \ln n}[1+\ln \ln \ln n-\ln 3-\ln \ln n]} \\
& \leq e^{\frac{3 \ln n}{\ln \ln n}[\ln \ln \ln n-\ln \ln n]} \\
& =e^{\left[-3 \ln n+\frac{3 \ln n \ln \ln \ln n}{\ln \ln n}\right]}
\end{aligned}
$$

For large values of $n,-3 \ln n+\frac{3 \ln n \ln \ln \ln n}{\ln \ln n} \leq-2 \ln n$.
Thus, $\left(\frac{e}{k}\right)^{k} \leq e^{-2 \ln n}=\frac{1}{n^{2}}$

## References

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2. Introduction to Probability by Blitzstein and Hwang, CRC Press 2015.
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4. My Notes on Algorithm Design.
