

Quick Review of Probability

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Sample Space & Events

Random Variable

Geometric Distribution

Coupon Collector Problem

Sample Space & Events

Basic Definition

Definitions

Sample Space S = Set of Outcomes.

Events \mathcal{E} = Subsets of S .

Probability is a function from subsets $A \subseteq S$ to positive real numbers between $[0, 1]$ such that:

1. $Pr(S) = 1$
2. For all $A, B \subseteq S$ if $A \cap B = \emptyset$, $Pr(A \cup B) = Pr(A) + Pr(B)$.
3. If $A \subset B \subseteq S$, $Pr(A) \leq Pr(B)$.
4. Probability of complement of A , $Pr(\bar{A}) = 1 - Pr(A)$.

Basic Definition

Examples:

1. Flipping a fair coin:

$$S = \{H, T\};$$

$$\mathcal{E} = \{\emptyset, \{H\}, \{T\}, S = \{H, T\}\}$$

2. Flipping fair coin twice:

$$S = \{HH, HT, TH, TT\};$$

$$\mathcal{E} = \{\emptyset, \{HH\}, \{HT\}, \{TH\}, \{TT\},$$

$$\{HH, TT\}, \{HH, TH\}, \{HH, HT\},$$

$$\{HT, TH\}, \{HT, TT\}, \{TH, TT\},$$

$$\{HH, HT, TH\}, \{HH, HT, TT\}, \{HH, TH, TT\},$$

$$\{HT, TH, TT\}, S = \{HH, HT, TH, TT\}\}$$

3. Rolling fair die twice:

$$S = \{(i, j) : 1 \leq i, j \leq 6\};$$

$$\mathcal{E} = \{\emptyset, \{1, 1\}, \{1, 2\}, \dots, S\}$$

Random Variable

Expectation

Definition

A random variable X is a function from sample space S to Real numbers, $X : S \rightarrow \mathfrak{R}$.

Expected value of a discrete random variable X is given by $E[X] = \sum_{s \in S} X(s) * Pr(X = X(s))$.

Note: Its a misnomer to say X is a random variable, it's a function.

Example: Flip a fair coin and define the random variable $X : \{H, T\} \rightarrow \mathfrak{R}$ as

$$X = \begin{cases} 1 & \text{Outcome is Heads} \\ 0 & \text{Outcome is Tails} \end{cases}$$

$$E[X] = \sum_{s \in \{H, T\}} X(s) * Pr(X = X(s)) = 1 * \frac{1}{2} + 0 * \frac{1}{2} = \frac{1}{2}$$

Linearity of Expectation

Definition

Consider two random variables X, Y such that $X, Y : S \rightarrow \mathfrak{R}$, then $E[X + Y] = E[X] + E[Y]$.

In general, consider n random variables X_1, X_2, \dots, X_n such that $X_i : S \rightarrow \mathfrak{R}$, then $E[\sum_{i=1}^n X_i] = \sum_{i=1}^n E[X_i]$.

Example: Flip a fair coin n times and define n random variable X_1, \dots, X_n as

$$X_i = \begin{cases} 1 & \text{Outcome is Heads} \\ 0 & \text{Outcome is Tails} \end{cases}$$

$$E[X_1 + \dots + X_n] = E[X_1] + \dots + E[X_n] = \frac{1}{2} + \dots + \frac{1}{2} = \frac{n}{2}$$

= Expected # of Heads in n tosses.

Proof of Linearity of Expectation

$$\begin{aligned}E[X + Y] &= \sum_{\omega \in S} (X + Y)[\omega] \cdot P(\omega) \\&= \sum_{\omega \in S} (X[\omega] + Y[\omega]) \cdot P(\omega) \\&= \sum_{\omega \in S} (X[\omega] \cdot P(\omega) + Y[\omega] \cdot P(\omega)) \\&= \sum_{\omega \in S} X[\omega] \cdot P(\omega) + \sum_{\omega \in S} Y[\omega] \cdot P(\omega) \\&= E[X] + E[Y]\end{aligned}$$

This generalizes to the sum of n random variables:

$$E[X_1 + \cdots + X_n] = E[X_1] + \cdots + E[X_n].$$

More detailed proof

$$\begin{aligned}E[X + Y] &= \sum_x \sum_y (x + y)P(X = x, Y = y) \\&= \sum_x \sum_y xP(X = x, Y = y) + \sum_x \sum_y yP(X = x, Y = y) \\&= \sum_x \sum_y xP(X = x, Y = y) + \sum_y \sum_x yP(X = x, Y = y) \\&= \sum_x x \sum_y P(X = x, Y = y) + \sum_y y \sum_x P(X = x, Y = y) \\&= \sum_x xP(X = x) + \sum_y yP(Y = y) \\&= E[X] + E[Y]\end{aligned}$$

Geometric Distribution

Geometric Distribution

Definition

Perform a sequence of independent trials till the first success. Each trial succeeds with probability p (and fails with probability $1 - p$).

A Geometric Random Variable X with parameter p is defined to be equal to $n \in \mathbb{N}$ if the first $n - 1$ trials are failures and the n -th trial is success. Probability distribution function of X is

$$Pr(X = n) = (1 - p)^{n-1}p.$$

Let Z to be the *r.v.* that equals the # failures before the first success, i.e. $Z = X - 1$.

Problem: Evaluate $E[X]$ and $E[Z]$.

To show: $E[Z] = \frac{1-p}{p}$ and $E[X] = 1 + \frac{1-p}{p} = \frac{1}{p}$.

Computation of $E[Z]$

$Z = \#$ failures before the first success.

Set $q = 1 - p$.

- $Pr(Z = k) = q^k p$
- $\frac{1}{1-q} = \sum_{k=0}^{\infty} q^k$ (for $0 < q < 1$)
- $\frac{1}{(1-q)^2} = \sum_{k=0}^{\infty} kq^{k-1}$

$$\begin{aligned} E[Z] &= \sum_{k=0}^{\infty} k Pr(Z = k) \\ &= \sum_{k=0}^{\infty} kq^k p \\ &= pq \sum_{k=0}^{\infty} kq^{k-1} \\ &= \frac{pq}{(1-q)^2} = \frac{1-p}{p} \end{aligned}$$

Examples

Examples:

1. Flipping a fair coin till we get a Head:

$$p = \frac{1}{2} \text{ and } E[X] = \frac{1}{p} = 2$$

2. Roll a die till we see a 6:

$$p = \frac{1}{6} \text{ and } E[X] = \frac{1}{p} = 6$$

3. Keep buying LottoMax tickets till we win (assuming we have 1 in 33294800 chance).

$$p = \frac{1}{33294800} \text{ and } E[X] = \frac{1}{p} = 33,294,800.$$

Coupon Collector Problem

Coupon's Collector Problem

Problem Definition

There are a total of n different types of coupons (Pokemon cards). A cereal manufacturer has ensured that each cereal box contains a coupon. Probability that a box contains any particular type of coupon is $\frac{1}{n}$. What is the expected number of boxes we need to buy to collect all the n coupons?

Define r.v. N_1, N_2, \dots, N_n , where N_i = # of boxes bought till the i -th coupon is collected.

Each N_i is a geometric random variable.

Coupon's Collector Problem Contd.

Let $N = \sum_{j=1}^n N_j$; Note $N_1 = 1$

$$E[N_j] = \frac{1}{\text{Pr of success in finding the } j^{\text{th}} \text{ coupon}} = \frac{1}{\frac{n-j+1}{n}}$$

$E[N] = \sum_{j=1}^n \frac{n}{n-j+1} = nH_n$, where $H_n = n$ -th Harmonic Number.

$$H_n = \sum_{i=1}^n \frac{1}{i} \text{ and } \ln n \leq H_n \leq \ln n + 1.$$

Thus, $E[N] = nH_n \approx n \ln n$,

Is $E[N] = nH_n = n \ln n$ a good estimate?

What is the probability that $E[N]$ exceeds $2nH_n$? Applying Markov's

Inequality: $Pr(X > s) \leq \frac{E[X]}{s} Pr(N > 2nH_n) < \frac{E[N]}{2nH_n} = \frac{nH_n}{2nH_n} = \frac{1}{2}$

Can we have a better bound?

Next: We show $Pr(N > n \ln n + cn) < \frac{1}{e^c}$

Pr. of missing a coupon after $n \ln n + cn$ boxes have been bought

$$= \left(1 - \frac{1}{n}\right)^{n \ln n + cn} \leq e^{-\frac{1}{n}(n \ln n + cn)} = \frac{1}{ne^c}.$$

Pr. of missing at least one coupon $\leq n\left(\frac{1}{ne^c}\right) = \frac{1}{e^c}$.

References

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