## Assignment 1

## COMP 3804, Fall 2021

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## 1 Guidelines

General guidelines are as follows:

- 1. Since we are only accepting assignments via the Brighspace system, no late submissions will be entertained after the cut-off time & date.
- 2. Please write clearly and answer questions precisely. It is your responsibility to ensure that what is uploaded is clearly readable. If we can't read, we can't mark!
- 3. Please cite all the references (including web-sites, names of friends, etc.) which you used/consulted as the source of information for each of the questions.
- 4. The first ten questions are worth 10 points each, totalling 100. The maximum mark for this assignment is 100. Bonus problem has ten extra points.
- 5. When a question asks you to design an algorithm it requires you to
  - (a) Clearly spell out the **steps** of your algorithm in pseudo code.
  - (b) **Prove** that your algorithm is correct
  - (c) **Analyze** the running time.
- 6. This assignment is based on the material that should have been learnt in COMP 1805/2402/2804. The purpose is to ensure that you have a good grasp of the background material to understand various topics in this course. If you have difficulty answering any of these questions, it is your responsibility to review them quickly.

## 2 Problems

- 1. (a) Define the functions O(),  $\Omega()$  and  $\Theta()$  that occur frequently in analyzing the complexity of algorithms (refer to your notes/textbooks of COMP 1805/2402/2804).
  - (b) Let  $p(n) = a_d n^d + a_{d-1} n^{d-1} + \cdots + a_1 n + a_0$ , where  $a_d > 0$ , be a d-degree polynomial in n. Also  $a_0, \dots, a_d$  are positive constants. Let k be a positive integer. Using your definitions in Part (a) show that:
    - i. If k > d, then  $p(n) \in O(n^k)$ .
    - ii. If k < d, then  $p(n) \in \Omega(n^k)$ .
    - iii. Is  $20n^3 + 10n^2 100n 5 \in O(n^4)$ ?
    - iv. Is  $20n^3 + 10n^2 100n 5 \in \Omega(n^2)$ ?
- 2. Evaluate the following recurrences (You can assume that T(1) = 1 for each of them).
  - (a) T(n) = 5T(n/5) + O(n)
  - (b) T(n) = T(8n/9) + O(n)
  - (c) T(n) = T(n-2) + O(1)
  - (d)  $T(n) = \sqrt{n}T(\sqrt{n}) + n$
- 3. Suppose you need to choose between the following algorithms which solves the same problem:
  - (a) Algorithm A solves the problem by dividing it into 4 subproblems of half of the size, recursively solves each of them, and combines the solution in linear time.
  - (b) Algorithm B solves the problem of size n by recursively solving three subproblems of size n-6 and then combining the solutions in constant time.
  - (c) Algorithm C solves the problem of size n by dividing it into 9 subproblems of size n/3 each, recursively solving each of them, and then combining the solution in  $O(n^2)$  time.

What are the running times of each of these algorithms? Which one will you choose and Why?

- 4. We want to sort n > 0 distinct real numbers in ascending order. Assume that these numbers are given in an array A of size n. We are also given a function double-partition(i,j) which takes as input two indices  $1 \le i < j \le n$  of A, where  $j-i \ge 2$ , and returns two elements  $x, y \in \{A[i], A[i+1], \ldots, A[j]\}$  that satisfy the following:
  - (a) x < y
  - (b) The number of elements in  $\{A[i], A[i+1], \ldots, A[j]\}$  that are smaller than x are at most  $\lceil \frac{j-i}{3} \rceil$ .

- (c) The number of elements in  $\{A[i], A[i+1], \ldots, A[j]\}$  that are larger than y are at most  $\lceil \frac{j-i}{3} \rceil$ .
- (d) The number of elements in  $\{A[i], A[i+1], \ldots, A[j]\}$  that are larger than x but smaller than y are at most  $\lceil \frac{j-i}{3} \rceil$ .
- (e) It takes O(j-i) time to compute x and y.

Design an algorithm, running in  $O(n \log n)$  time, to sort any set of n distinct real numbers using the function double-partition. (Please review Guideline #5 in the context of what is expected when you are asked to design an algorithm.)

- 5. You are given an array A consisting of n positive integers, where each element is  $\leq 1000n$ . Devise an algorithm, running in O(n) time, to sort A in ascending order. Justify your answer.
- 6. Recall (from your COMP 2402 course) that there is a lower bound for sorting. Answer the following questions
  - (a) What does it mean to have a lower bound for a problem?
  - (b) State clearly what is the lower bound claim for sorting a set of n (real)-numbers.
  - (c) Why this claim does not apply to the problem in Question 5?
- 7. You are given an array A consisting of n real numbers. Describe and analyze an algorithm, running in O(n) time, that rearranges the elements of A so that A forms a binary heap. Once A is transformed into a Binary Heap, show how you can report the elements in A in sorted (ascending) order. How much time it takes to report all the elements of A in the sorted order? Justify your answer.
- 8. You are given a set of n real numbers which you are asked to insert incrementally in an initially empty binary search tree. Note that the time to insert an element in a binary search tree of size x is  $O(\log x)$ . What is the total running time of inserting all the n elements in the tree. Justify your answer. Assume that you have formed the binary search tree on n elements, show how you can report the elements in a sorted order in O(n) time.
- 9. Reflecting on the answers to the previous two questions, is the construction of a binary heap in O(n) time or reporting the elements in sorted order from a binary search tree in O(n) time is in contradiction to the lower bound for sorting (refer to Question 6)? Justify your answer.
- 10. Discuss a couple of scenarios where you will be using a binary heap instead of a binary search tree. Justify your answer.
- 11. (Bonus Problem:) Assume that you have a set of three parallel and distinct lines in plane. Assume that Line 1 contains a set R of n red points, Line 2 contains a set R

of n blue points, and Line 3 contains a set G of n green points, where n is a positive integer. You need to design an algorithm that determines if there is a triplet of points (r, b, g) such that they lie on a line segment, where  $r \in R$ ,  $b \in B$  and  $g \in G$ . (Observe that it is straightforward to design an algorithm running in  $O(n^3)$  time by considering each choice for points r, b, and g, and in O(1) time testing whether they all lie on a segment.) To get Bonus points, design an algorithm running in  $O(n^2)$  time.

