

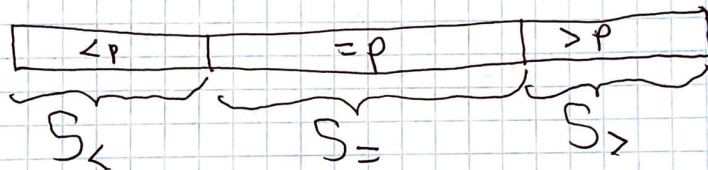
Randomized Selection

Algorithm $RSelect(S, k)$:

// S is a sequence of numbers, $1 \leq k \leq |S|$,
 // returns the k -th smallest element in S

if $|S| = 1$: return the only element in S

if $|S| \geq 2$: $p =$ uniformly random element in S ;
 by scanning S , divide it into



if $k \leq |S_<|$: $RSelect(S_<, k)$

if $|S_<| + 1 \leq k \leq |S_<| + |S_|=|$: return p

if $k \geq |S_<| + |S_|=| + 1$:

$RSelect(S_>, k - |S_<| - |S_|=|)$

Let n be a large integer, let S be a sequence

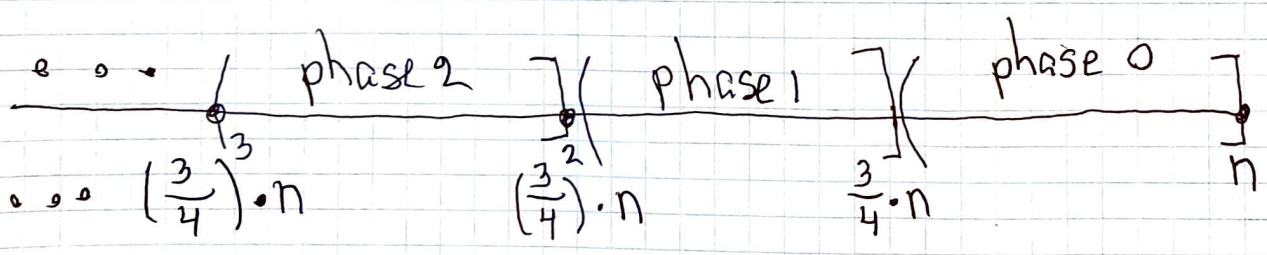
of length n . Define random variable

$T =$ time when running $RSelect(S, k)$.

We are going to show: $E(T) = O(n)$.

When we run $RSelect(S, k)$, recursive calls are generated on smaller and smaller sequences.

For $i = 0, 1, 2, \dots$: a call is in phase i if the length of the sequence in this call is $\leq (\frac{3}{4})^i \cdot n$ and $> (\frac{3}{4})^{i+1} \cdot n$



Initially: phase 0, because $|S| = n$

During recursive calls, we either stay in the

same phase or move to a phase with a larger index.

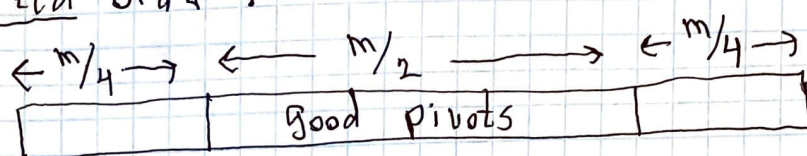
Which case happens depends on the pivot p (which is randomly chosen).

Consider a call in phase i .

$m =$ length of the sequence in this call.

$$\left(\frac{3}{4}\right)^{i+1} \cdot n < m \leq \left(\frac{3}{4}\right)^i \cdot n$$

For purpose of analysis, consider these m numbers in sorted order:



pivot p is good if p is in the middle half

$$* \Pr(\text{pivot } p \text{ is good}) = \frac{1}{2}$$

* Assume pivot p is good. If there is a next call, then it is on a sequence of length

$$\leq m - \frac{m}{4} = \frac{3}{4} \cdot m \leq \left(\frac{3}{4}\right)^{i+1} \cdot n$$

\therefore this next call is in a phase $\geq i+1$

From COMP 2804:

RS4

experiment $\begin{cases} \rightarrow \text{success with probability } \alpha \\ \rightarrow \text{failure with probability } 1-\alpha \end{cases}$

Repeat experiment until success ~~at~~ for the first time.

$$\text{Then: } E(\# \text{ times}) = 1/\alpha$$

Define random variable

$$X_i = \# \text{ calls in phase } i$$

Then: $E(X_i) \leq 2$ [Exercise: Why ≤ 2 ?
Why not $= 2$?]

* time for one call in phase i (excluding the recursive call)

$$\leq cm \leq c \left(\frac{3}{4}\right)^i \cdot n$$

for some constant c .

Remember: T = time when running

$R_{\text{select}}(S, k)$, where $|S| = n$.

$$T \leq \sum_{i \geq 0} c \left(\frac{3}{4}\right)^i \cdot n \cdot X_i$$

Using linearity of expectation:

$$E(T) \leq \sum_{i \geq 0} c \left(\frac{3}{4}\right)^i \cdot n \cdot E(X_i)$$

$$\leq 2cn \sum_{i=0}^{\infty} \left(\frac{3}{4}\right)^i$$

$$= 2cn \cdot \frac{1}{1 - 3/4}$$

$$= 8cn$$

$$= O(n)$$

$$\sum_{i=0}^{\infty} x^i = \frac{1}{1-x}$$