

Intelligent Game Playing

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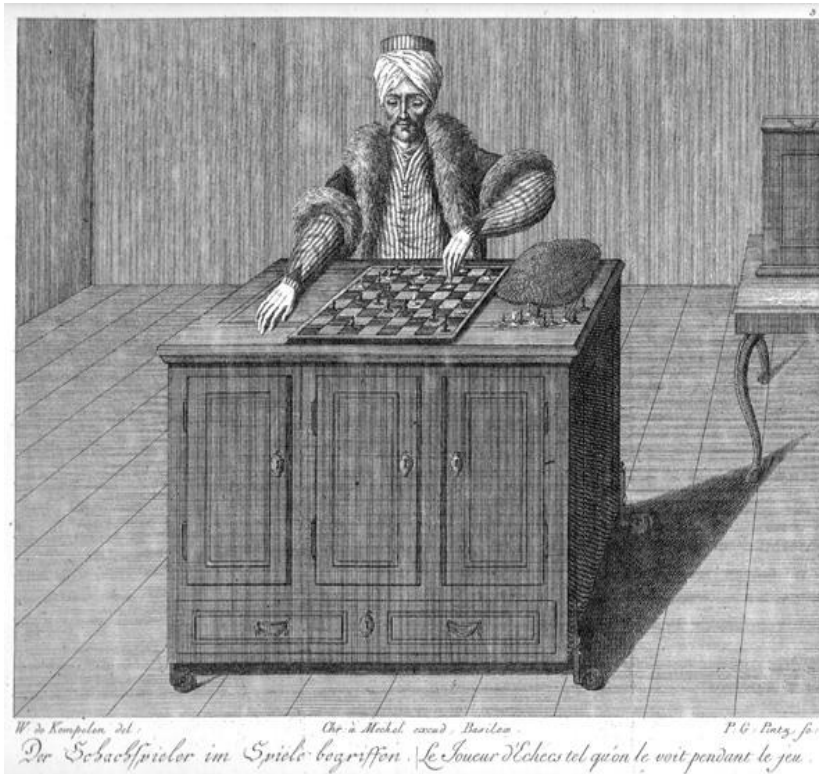
The primary source of these notes are the slides of Professor Hwee Tou Ng from Singapore. The Multi-Player Game section was due to Mr. Spencer Polk.

I sincerely thank them for this.

Games - Introduction

- Question: Can machines outplay humans?
- Captured imaginations for centuries
 - Appearance in myth and legend
 - Popular topic in fiction
- Thanks to AI and search techniques, the dream has come true!

History



The Turk (1770)

- “The Turk”: In 1770
- A chess playing machine
- Toured Europe
- Facing well-known opponents
 - e.g. Napoleon, Ben Franklin
- Of course: Revealed - fraud

History

- “The Turk” shows how fascinating this idea is
- 1914: King vs Rook strategies by automaton
- True AI game playing – Claude Shannon: 1950
 - Based on earlier work by Nash and Neumann
- Shannon's algorithm still used
 - **Mini-Max Search** (we will return to it shortly)

History

- Shannon's 1950 paper focused on Chess
 - Chess remains very important to game playing research
- At the time, seen as purely theoretical exercise
- 1970s: First commercial Chess programs
- 1980s: Chess programs playing at Expert level
 - Still some time until Grandmaster level...

History



Kasparov vs Deep Blue

- 1997: IBM's Deep Blue
 - Defeats Garry Kasparov
- First defeat of Grandmaster
- Field: Branched out since
 - Poker, Go: Now important games
 - IBM Watson on Jeopardy

Games vs. Search Problems

- “Unpredictable” opponent
 - Specifying a move for every possible opponent reply
- Time limits
 - Unlikely to find goal, must approximate

Mini-Max Search

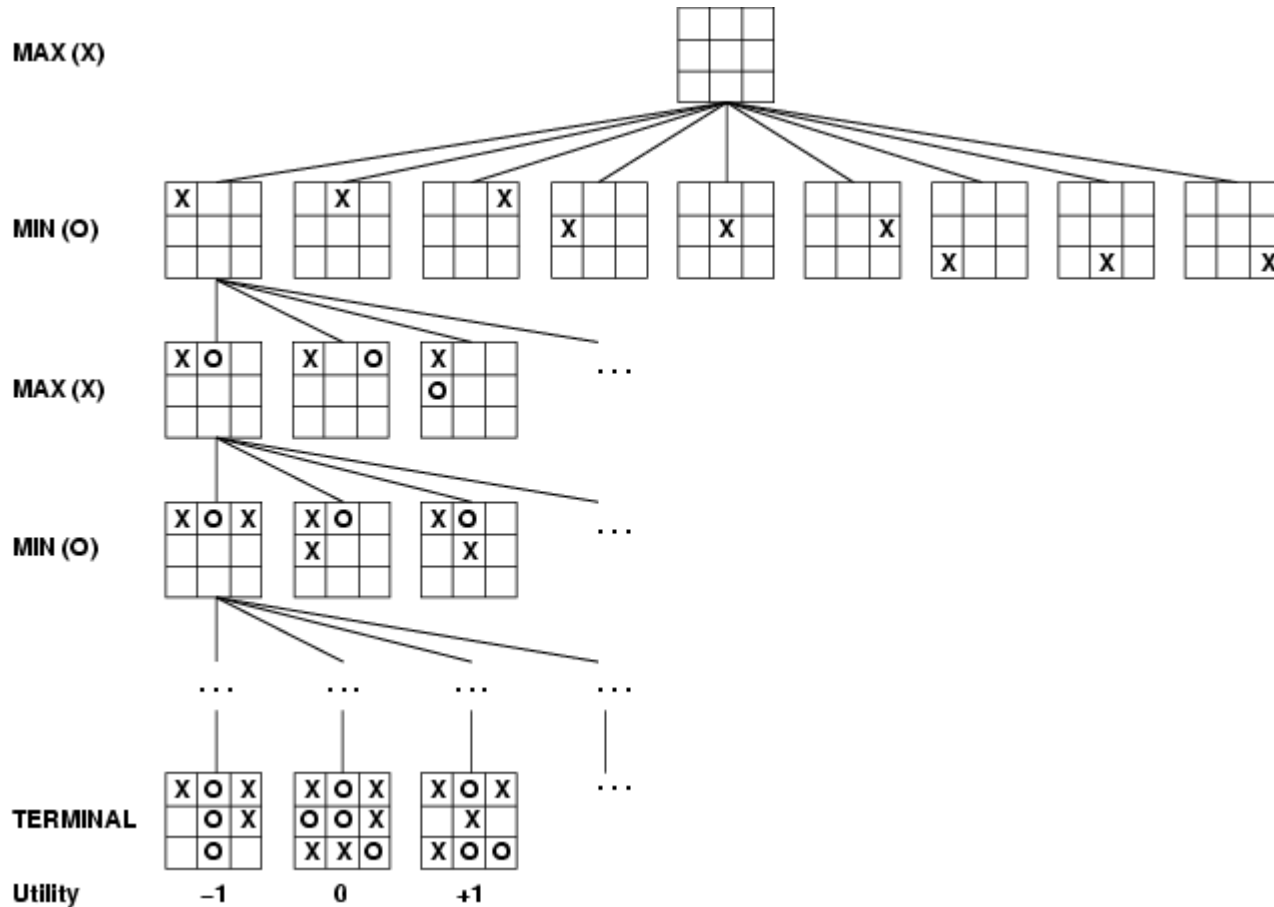
- Search to find the correct move in a two player game
- The **optimal** solution:
 - Exponential algorithm
 - Generate all possible paths
 - Only play those that lead to a winning final position
- **Realistic** alternative to the Optimal
- Use finite depth look-ahead with a heuristic function
- Evaluate how good a given game state is

Mini-Max

- Extend Tree down to a given search depth
- Top of tree is the **Computer's** move
 - Wants move to ultimately be one step closer to a winning position
 - Wants move that **maximizes** own chance of winning
- Next move is **Opponent's**
 - Opponent assumed to perform a move that his best
 - Wants move that **minimizes** Computer's chance of winning

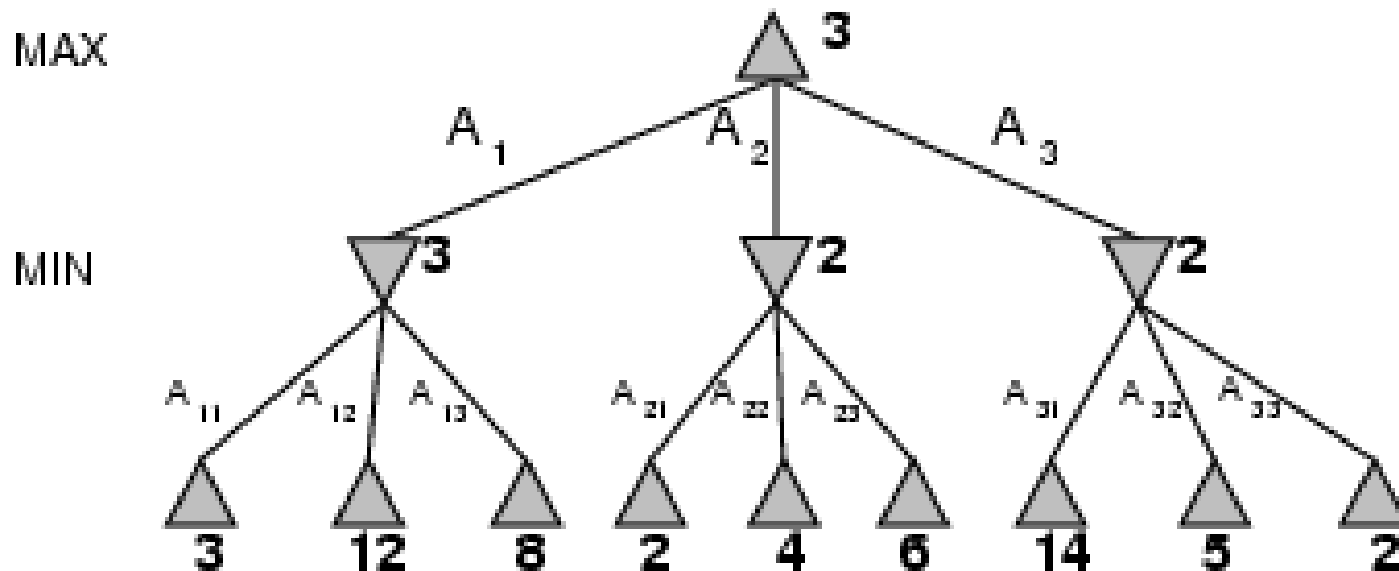
Game tree

2-player, Deterministic, Turns



Mini-Max

- **Perfect** play for deterministic games
- **Idea**: Choose move to position with highest **Mini-Max value**
= **Best achievable payoff against best play**
- Example: 2-ply game:



Mini-Max for Nim

- Nim Game

- Two players start with a pile of tokens
- Legal move: Split (any) existing pile into two non-empty differently sized piles
- Game ends when no pile can be unevenly split
- Player who cannot make his move loses the game

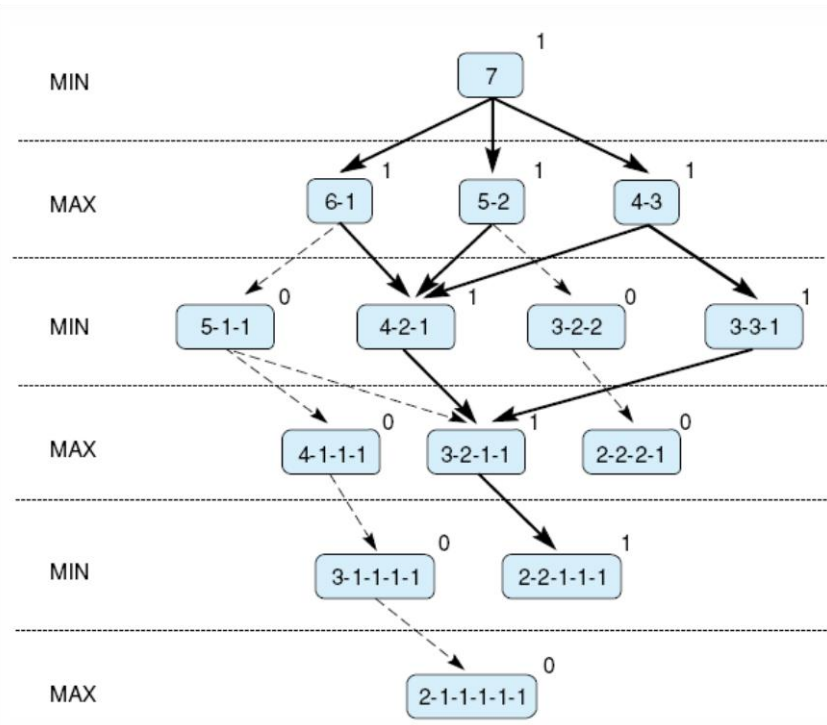
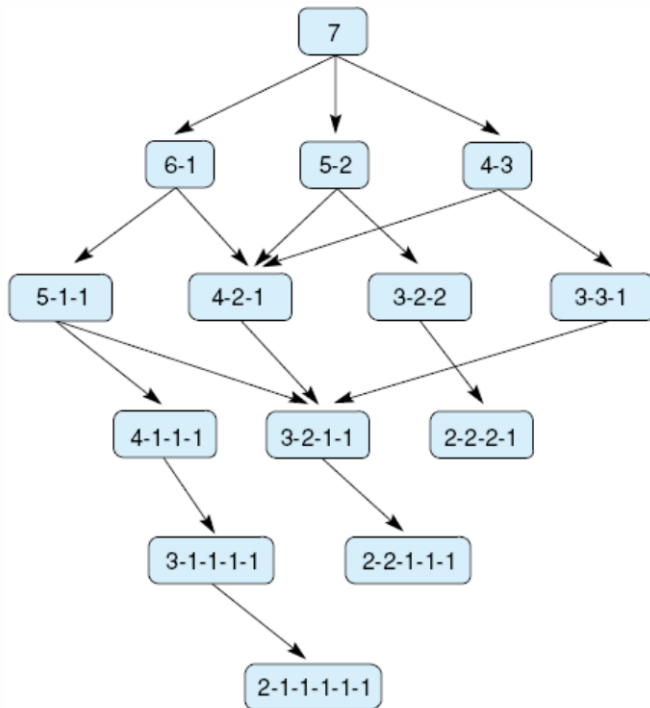
- Search strategy

- Existing heuristic search methods do not work

Mini-Max for Nim

- Label nodes as MIN or MAX, alternating for each level
- Define utility function (payoff function).
- Do **full** search on tree
 - Expand all nodes until game is over for each branch
- Label leaves according to outcome
- Propagate result up the tree with:
 - $M(n) = \max(\text{child nodes})$ for a MAX node
 - $m(n) = \min(\text{child nodes})$ for a MIN node
- Best next move for MAX is the one leading to the child with the highest value (and vice versa for MIN)

Mini-Max for Nim



Mini-Max Algorithm

```
function MINIMAX-DECISION(game) returns an operator
for each op in OPERATORS[game] do
    VALUE[op] := MIN-VALUE(APPLY(op, game), game)
end
return the op with the highest VALUE[op]
```

```
function MAX-VALUE(state, game) returns a utility value
if CUTOFF-TEST(state,) then return EVAL(state)
value :=  $-\infty$ 
for each s in SUCCESSORS(state) do
    value := MAX(value, MIN-VALUE(s, game))
end
return value
```

```
function MIN-VALUE(state, game) returns a utility value
if CUTOFF-TEST(state,) then return EVAL(state)
value :=  $\infty$ 
for each s in SUCCESSORS(state) do
    value := MIN(value, MAX-VALUE(s, game))
end
return value
```

Problems with Mini-Max

- **Horizon effect:** Can't see beyond depth
 - Due to exponential increase in tree size, only very limited depth feasible
 - Solution: Quiescence search. Start at the leaf nodes of the main search, and try to solve this problem.
 - In Chess, quiescence searches usually include all capture moves, so that tactical exchanges don't mess up the evaluation. In principle, quiescence searches should include any move which may destabilize the evaluation function--if there is such a move, the position is not quiescent.
- **May want to use look up tables**
 - For end games
 - Opening moves (called Book Moves)

Properties of Mini-Max

- **Complete?**
 - Yes (if tree is finite)
- **Optimal?**
 - Yes (against an optimal opponent)
- **Time complexity?**
 - $O(b^m)$
- **Space complexity?**
 - $O(bm)$ (depth-first exploration)
- **Chess:** $b \approx 35$, $m \approx 100$ for “reasonable” games
 - Exact solution completely infeasible

Branch and Bound: The α - β Algorithm

- **Branch and Bound**: If current path (branch) is already worse than some other known path:
 - Stop expanding it (bound).
- **Alpha-Beta** is a branch and bound technique for Mini-Max search
- If you know that the level above won't choose your branch because you have already found a value along one of your sub-branches that is too good, stop looking at other sub-branches that haven't been looked at yet

The α - β Algorithm

- Instead of maintaining a single mini-max value , the α - β pruning algorithm, maintains two: α , β
- Together provide a bound on the possible values of the mini-max tree at any given point.
- At any given point, α : **minimum** the player can expect to receive
- At any given point, β : **maximum** value the player can expect

The α - β Algorithm

- If it is ever the case that this bound is reversed or has range of 0 ($\beta \leq \alpha$), then better options exist for the player at other pre-explored nodes
- As α is the minimum value we know we can get
- Thus this node cannot be the mini-max value of the tree.
- There is no point in exploring any more of this node's children
- Potentially saving considerable computation time in a game with a large branching factor/depth

Properties of α - β

- Pruning **does not** affect final result
- Good move ordering improves pruning effectiveness
- With “perfect ordering” time complexity = $O(b^{m/2})$
 - **Doubles** depth of search
- α - β is a simple example of the value of reasoning about which computations are **really** relevant

Why it is called α - β

- α : Value of the best choice found so far at any choice point along the path for *max*
- If v is worse than α
 - *max* will avoid it
 - prune that branch
- Define β similarly for *min*

MAX

MIN

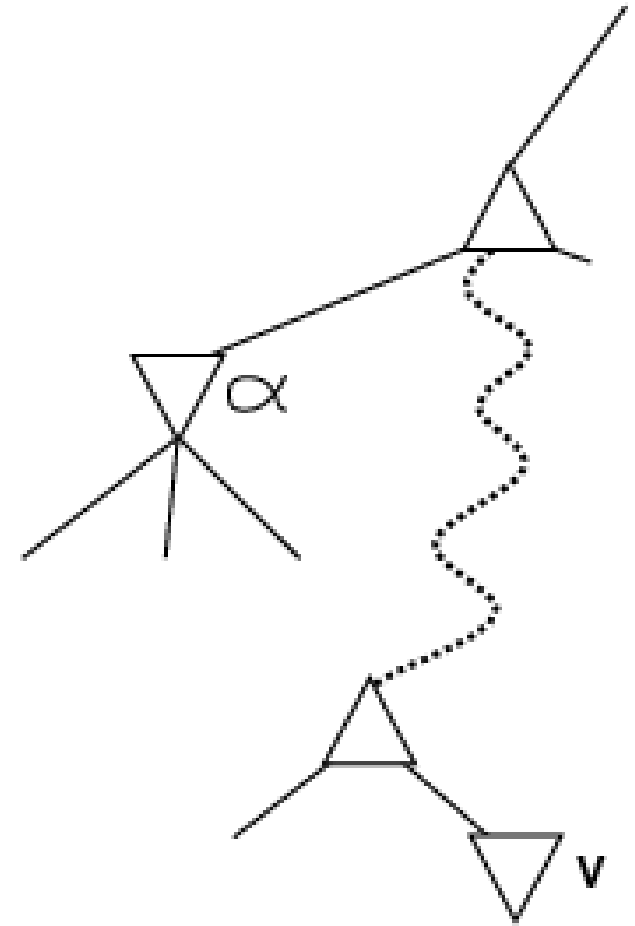
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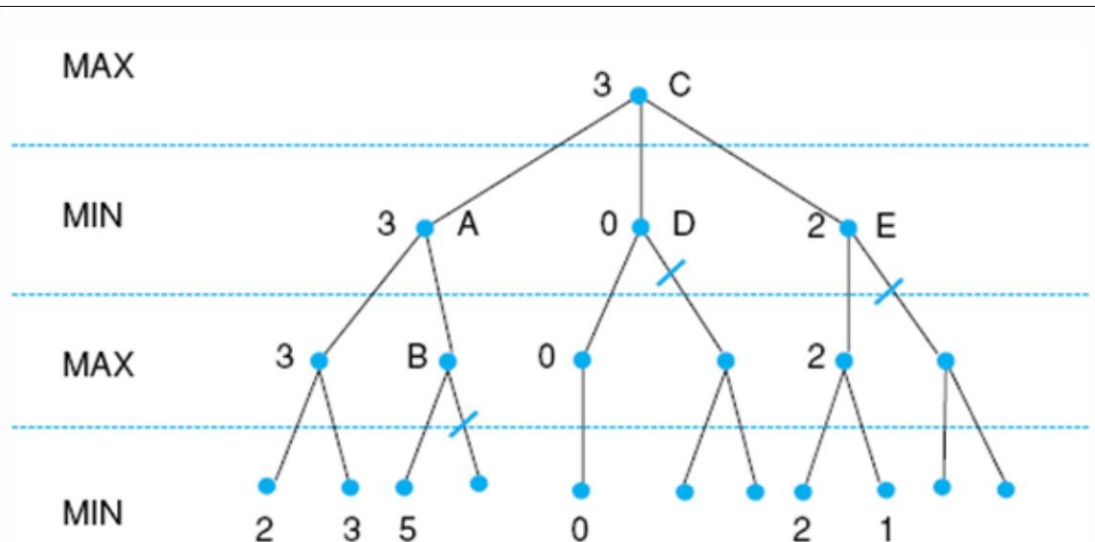
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MAX

MIN



Effects of α - β



A has $\beta = 3$ (A will be no larger than 3)

B is β pruned, since $5 > 3$

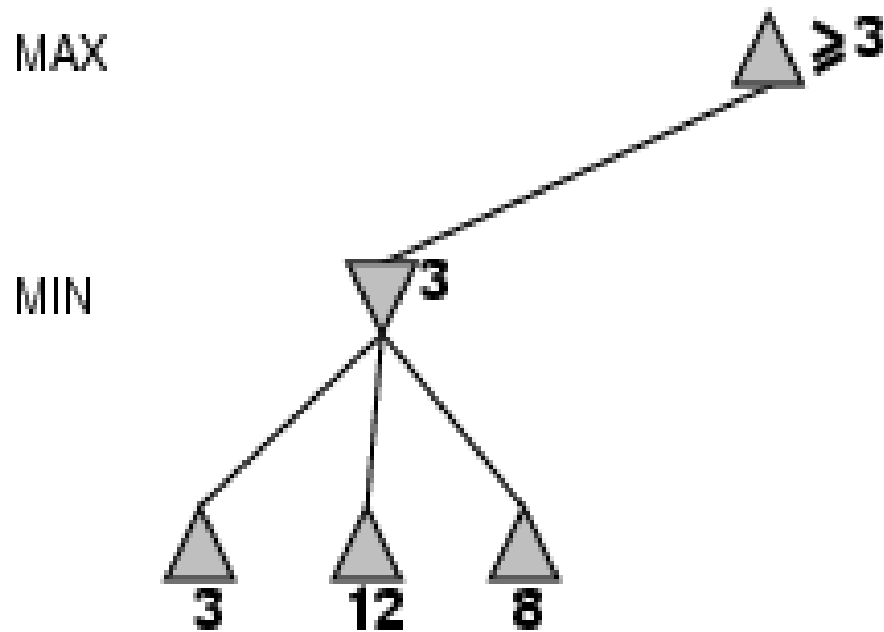
C has $\alpha = 3$ (C will be no smaller than 3)

D is α pruned, since $0 < 3$

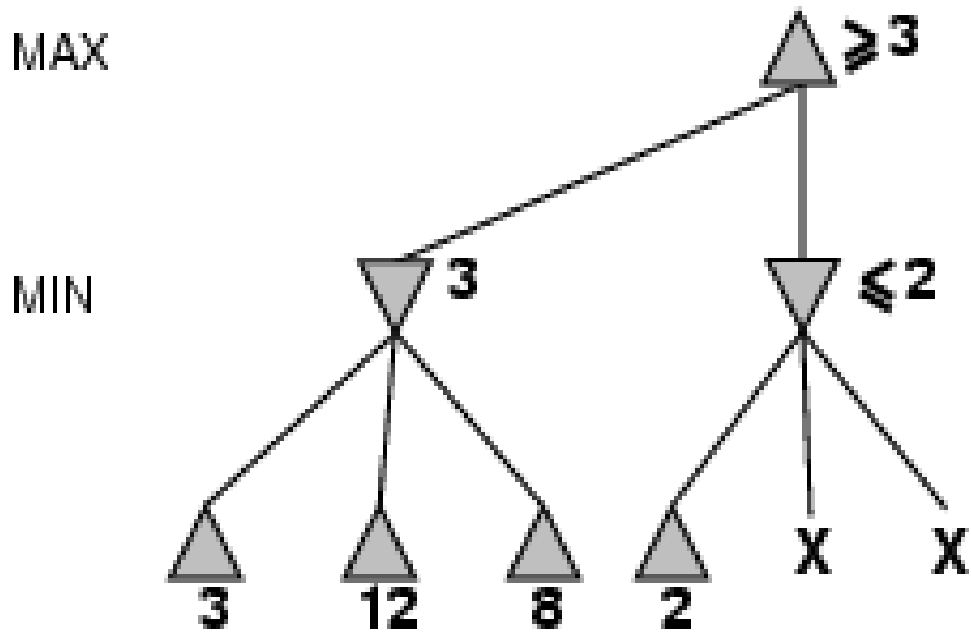
E is α pruned, since $2 < 3$

C is 3

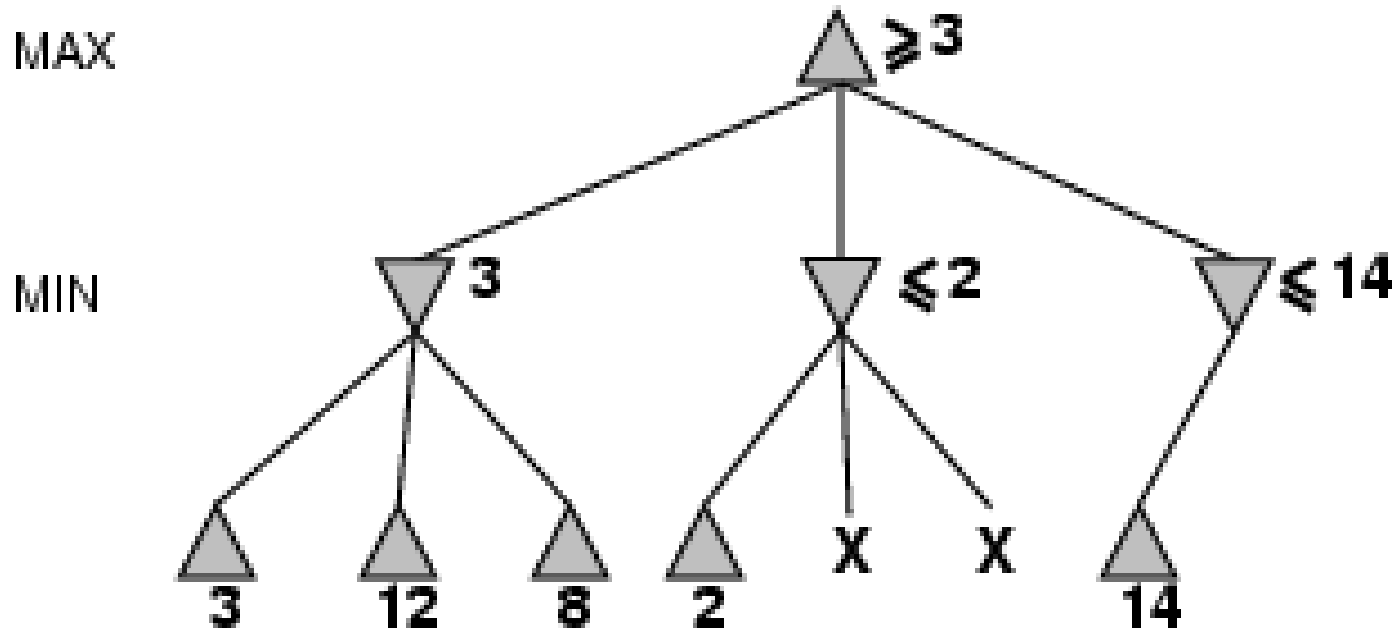
Example: α - β Pruning



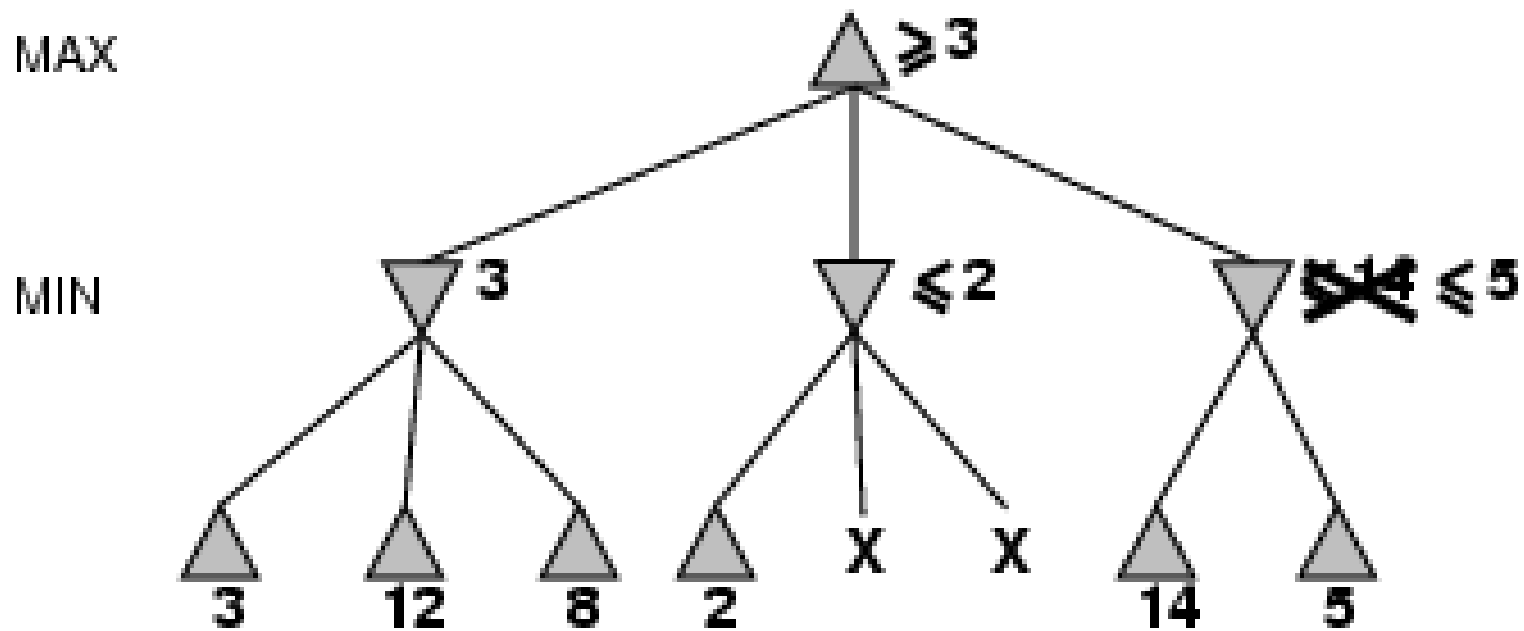
Example: α - β Pruning



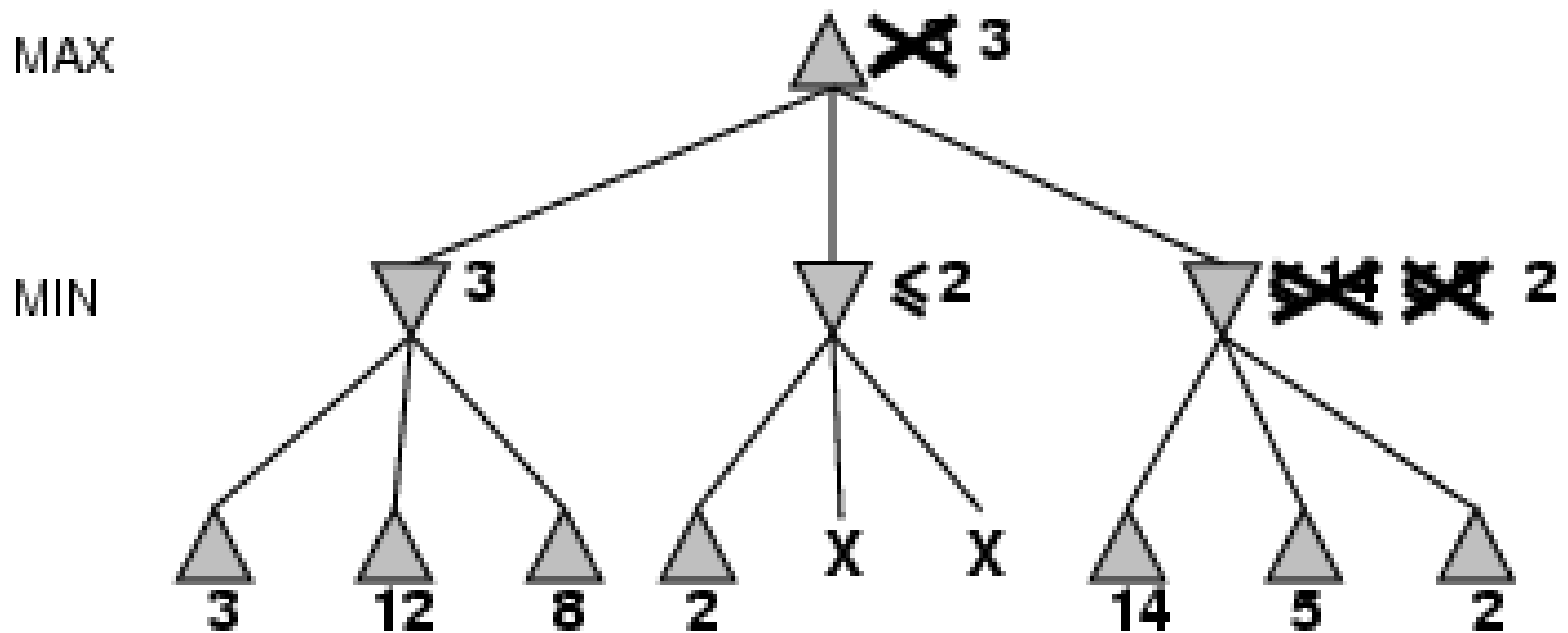
Example: α - β Pruning



Example: α - β Pruning



Example: α - β Pruning



The α - β Algorithm

- From Russell and Norvig α = best score for MAX so far game = game description
 β = best score for MIN so far state = current state in game

```
function MAX-VALUE(state, game,  $\alpha$ ,  $\beta$ ) returns a utility value  
if CUTOFF-TEST(state,) then return EVAL(state)  
for each s in SUCCESSORS(state) do  
     $\alpha :=$  MAX( $\alpha$ , MIN-VALUE(s, game,  $\alpha$ ,  $\beta$ ))  
    if  $\alpha \geq \beta$  then return  $\alpha$   
end  
return  $\alpha$ 
```

```
function MIN-VALUE(state, game,  $\alpha$ ,  $\beta$ ) returns a utility value  
if CUTOFF-TEST(state,) then return EVAL(state)  
for each s in SUCCESSORS(state) do  
     $\beta :=$  MIN( $\beta$ , MAX-VALUE(s, game,  $\alpha$ ,  $\beta$ ))  
    if  $\beta \leq \alpha$  then return  $\beta$   
end  
return  $\beta$ 
```

The α - β Algorithm

function ALPHA-BETA-SEARCH(*state*) *returns an action*

inputs: *state*, current state in game

$v \leftarrow \text{MAX-VALUE}(state, -\infty, +\infty)$

return the *action* in SUCCESSORS(*state*) with value v

function MAX-VALUE(*state*, α , β) *returns a utility value*

inputs: *state*, current state in game

α , the value of the best alternative for MAX along the path to *state*

β , the value of the best alternative for MIN along the path to *state*

if TERMINAL-TEST(*state*) **then return** UTILITY(*state*)

$v \leftarrow -\infty$

for a, s in SUCCESSORS(*state*) **do**

$v \leftarrow \text{MAX}(v, \text{MIN-VALUE}(s, \alpha, \beta))$

if $v \geq \beta$ **then return** v

$\alpha \leftarrow \text{MAX}(\alpha, v)$

return v

The α - β Algorithm

function MIN-VALUE(*state*, α , β) *returns a utility value*

inputs: *state*, current state in game

α , the value of the best alternative for MAX along the path to *state*

β , the value of the best alternative for MIN along the path to *state*

if TERMINAL-TEST(*state*) **then return** UTILITY(*state*)

$v \leftarrow +\infty$

for a, s in SUCCESSORS(*state*) **do**

$v \leftarrow \text{MIN}(v, \text{MAX-VALUE}(s, \alpha, \beta))$

if $v \leq \alpha$ **then return** v

$\beta \leftarrow \text{MIN}(\beta, v)$

return v

Improving Game Playing

- Increase Depth of Search
- Have better heuristic for game state evaluation

Changing Levels of Difficulty

- Increase Depth of Search

Resource Limits

- Suppose we have 100 secs, explore 10^4 nodes/sec
 - 10^6 nodes per move
- Standard approach:
 - **Cutoff test**: Depth limit (perhaps add **quiescence search**)
- Evaluation function:
 - Estimated desirability of position

Evaluation Functions

- Chess, typically **linear** weighted sum of **features**

$$Eval(s) = w_1 f_1(s) + w_2 f_2(s) + \dots + w_n f_n(s)$$

- Example: $w_1 = 9$ with
 $f_1(s) = (\text{number of white queens}) - (\text{number of black queens})$
etc.

Cutting-Off Search

MinimaxCutoff is identical to *MinimaxValue* except

1. *Terminal?* is replaced by *Cutoff?*
2. *Utility* is replaced by *Eval*

Does it work in practice?

$$b^m = 10^6, b=35 \rightarrow m=4$$

4-ply lookahead is a hopeless chess player!

- 4-ply \approx human novice
- 8-ply \approx typical PC, human master
- 12-ply \approx Deep Blue, Kasparov

Quiescence search

- Quiescence search: Study moves that are noisy
- They appear good, but moves around them - bad
- Investigate them with a localized leaf search
- Attempt to identify delaying tactics and change the seemingly-good value of the node
- A very natural extension of mini-max
- Simply run search again at a leaf node until that leaf node becomes quiet
- As with iterative deepening, running time of the algorithm won't increase by more than a constant

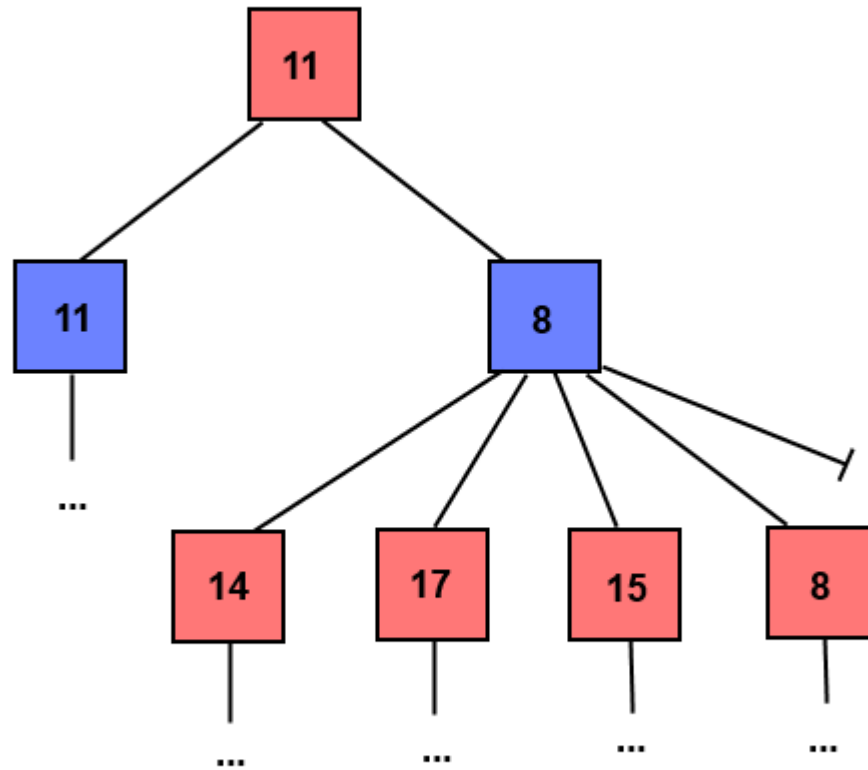
Real Deterministic Games

- **Checkers:** Chinook ended 40-year-reign of human world champion Marion Tinsley in 1994.
 - Used a precomputed endgame database
 - Defining perfect play for all positions involving 8 or fewer pieces on the board - a total of 444 billion positions.
- **Chess:** Deep Blue defeated human world champion Kasparov in a six-game match in 1997.
 - Deep Blue searches 200 million positions per second
 - Uses very sophisticated evaluation
 - Undisclosed methods for extending some lines of search up to 40 ply.

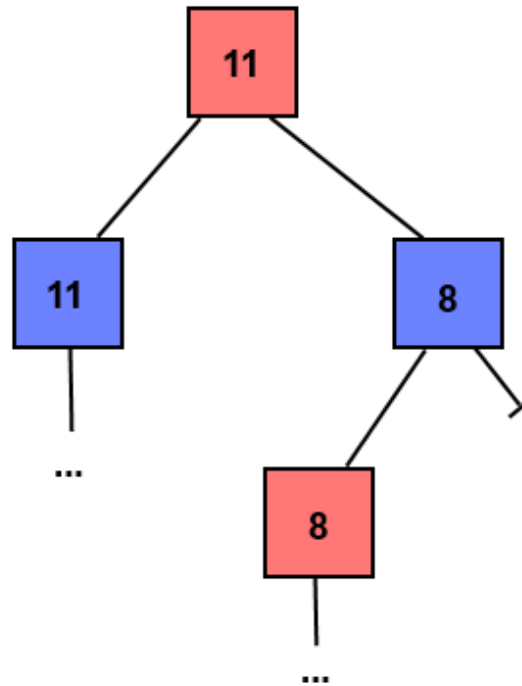
Move Ordering

- Best possible pruning is achieved if the best move is searched first at each level of the tree
- Problem: If we knew the best move, we would not need to search!
- Thus, we employ move ordering *heuristics*, which search the best move first
- **Example:** In Chess, search capturing moves before non-capturing moves
- **What we want:** domain *independent* techniques

Example: Poor Move Ordering



Example: Good Move Ordering



Principal Variation Move

- As it is a search algorithm, can apply Iterative Deepening to Mini-Max
- At each level, we thus find a move path we expect us and the opponent to take
- At the next stage, search it first!
 - Called **Principal Variation** move
- Even though Iterative Deepening takes some time, PV-move can greatly improve overall performance!

Other Heuristics

- **Killer Moves:** Remember move that produced a cut on this level of the tree
 - If we encounter it again, search it first!
 - Normally remember two moves per level
- **History Heuristic:** Same as Killer Moves, want to remember moves that produce cuts
 - Want to use info on all levels of tree
 - Hold array of counters, increment based on level cut occurred at
 - Details outside scope of this talk

Real Deterministic Games

- **Othello**: Human champions refuse to compete against computers, who are too good.

Things to Remember: Games

- Games are fun to work on!
- They illustrate several important points about AI
- Perfection is unattainable
- Must approximate paths and solutions
- Good idea to think about **what** to think about

Two Player to Multi-Player Games

- Mini-Max: Originally envisioned for Chess
 - Two player, deterministic, perfect information game
- What if we want to play a *multi*-player game?
 - Instead of two players, we have N players, where $N > 2$
 - Examples: Chinese Checkers, Poker
- New challenges, requiring new techniques!

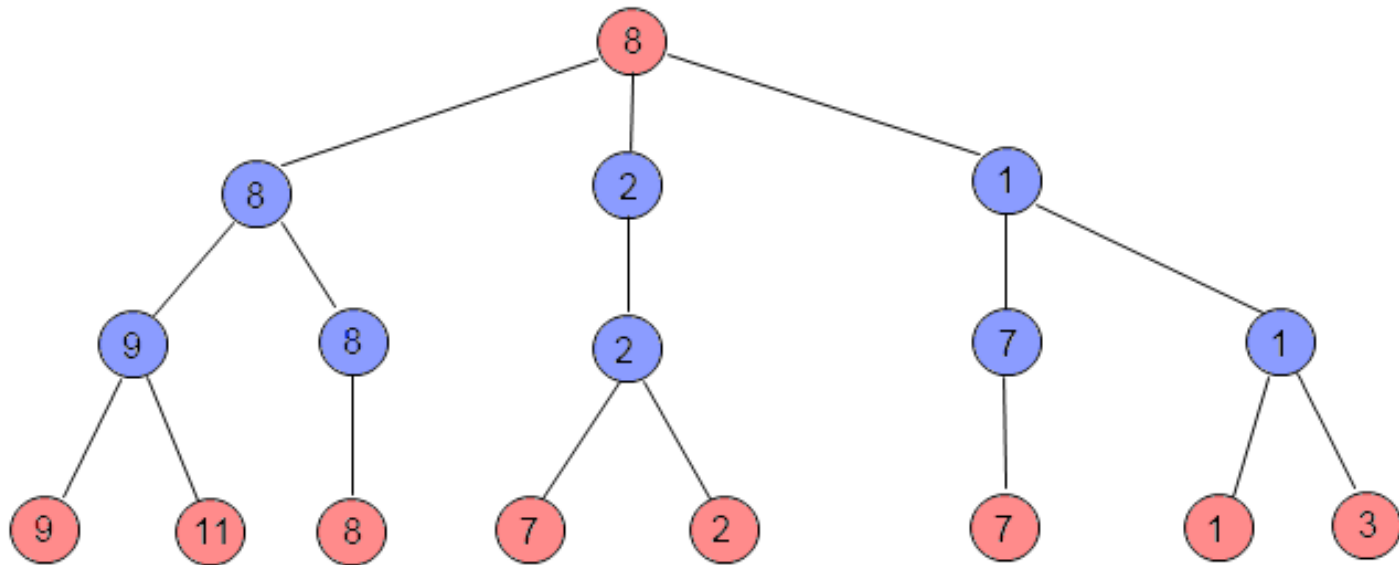
Qualities of Multi-Player Games

- In two player zero sum games, your gain is reflected in equal loss for opponent
 - No longer true for multi-player game
 - Loss spread between multiple opponents
- *Coalitions* may arise during play
- More opponent turns occur between perspectives

Extending Mini-Max to Multi-Player Games

- Problem: Mini-Max operates using a single value
 - Worked for two player games, as opponent's gain is our loss
- Single score very valuable – Allows **pruning**
 - Would like to keep pruning to speed up the search
- Simple solution: *All* opponents minimize our score
 - So, MAX-MIN-MIN, MAX-MIN-MIN-MIN, etc
- Called the *Paranoid Algorithm*

Paranoid Algorithm



Sample Paranoid Tree (Red MAX, Blue MIN)

Paranoid Algorithm

```
function integer paranoid(node, depth):
  if node is terminal or depth <= 0 then
    return heuristic value of node
  else
    if node is max then
      val =  $-\infty$ 
      for all child of node do
        val = max(val, paranoid(child, depth - 1))
      end for
    else
      val =  $\infty$ 
      for all child of node do
        val = min(val, paranoid(child, depth - 1))
      end for
    end if
  return val
end if
```

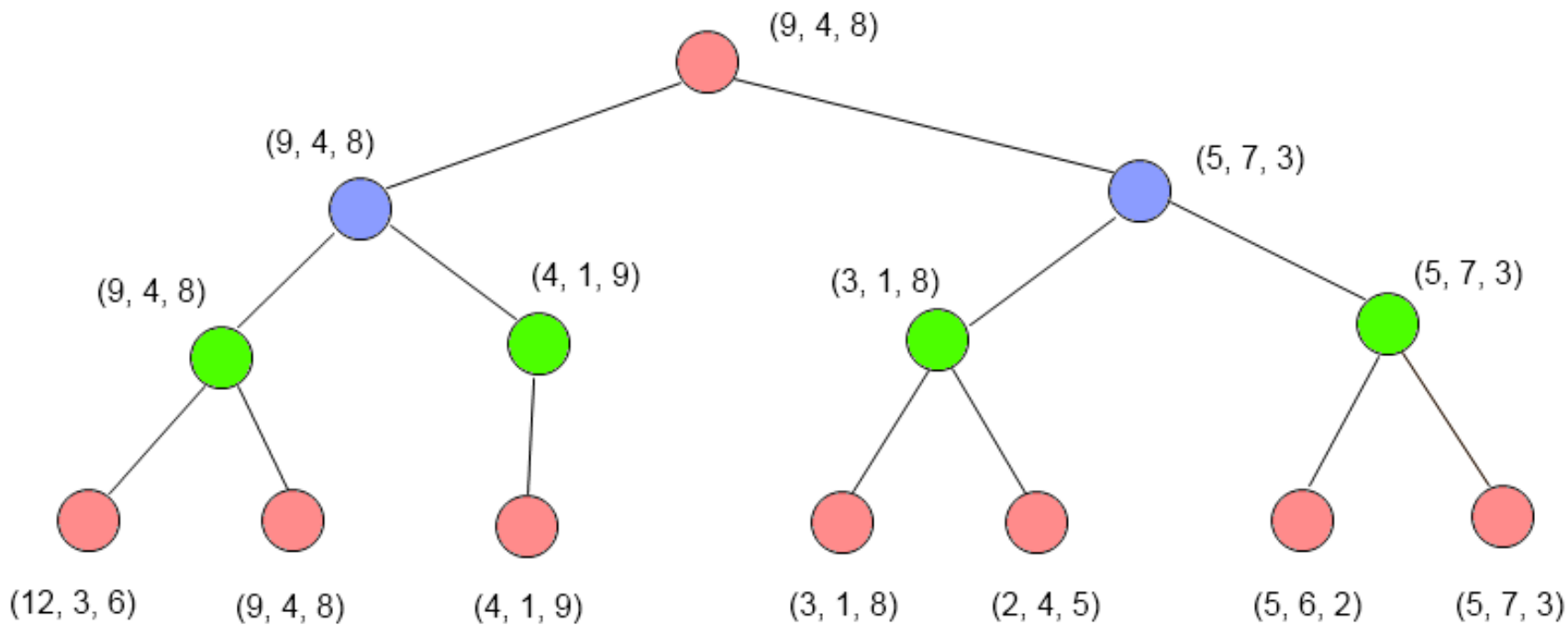
Paranoid Algorithm

- Algorithm *exact same* as Mini-Max in many implementations
- Pros
 - Easy to implement and understand
 - Subject to α - β pruning on MAX/MIN borders
 - Not for phases between MIN nodes
- Cons
 - Views all opponents as a coalition – leads to bad play
 - Limited look-ahead for perspective player
 - Need to have multiple MIN phases in a row

Max-N Algorithm

- 1986: Luckhardt and Irani
- Addresses coalition problem of Paranoid
- Keeps *tuple* of scores, not one value
- Assumption: Players maximize their *own* score
 - No consideration for other players
- Heuristic returns value for each player
 - i.e. [6, 3, 8] for three-player game
- *N*th player maximizes *N*th value

Max-N Algorithm



Sample Max-N Tree

Max-N Algorithm

```
function integer[] max-n(node, depth):
  if node is terminal or depth <= 0 then
    return heuristic value of node
  else
    val =  $-\infty$ 
    tuple = []
    for all child of node do
      val = max(val, max-n(child; depth - 1)[node.player])
      if val changed
        tuple = max-n(child; depth-1)
      end if
    end for
    return tuple
  end if
```

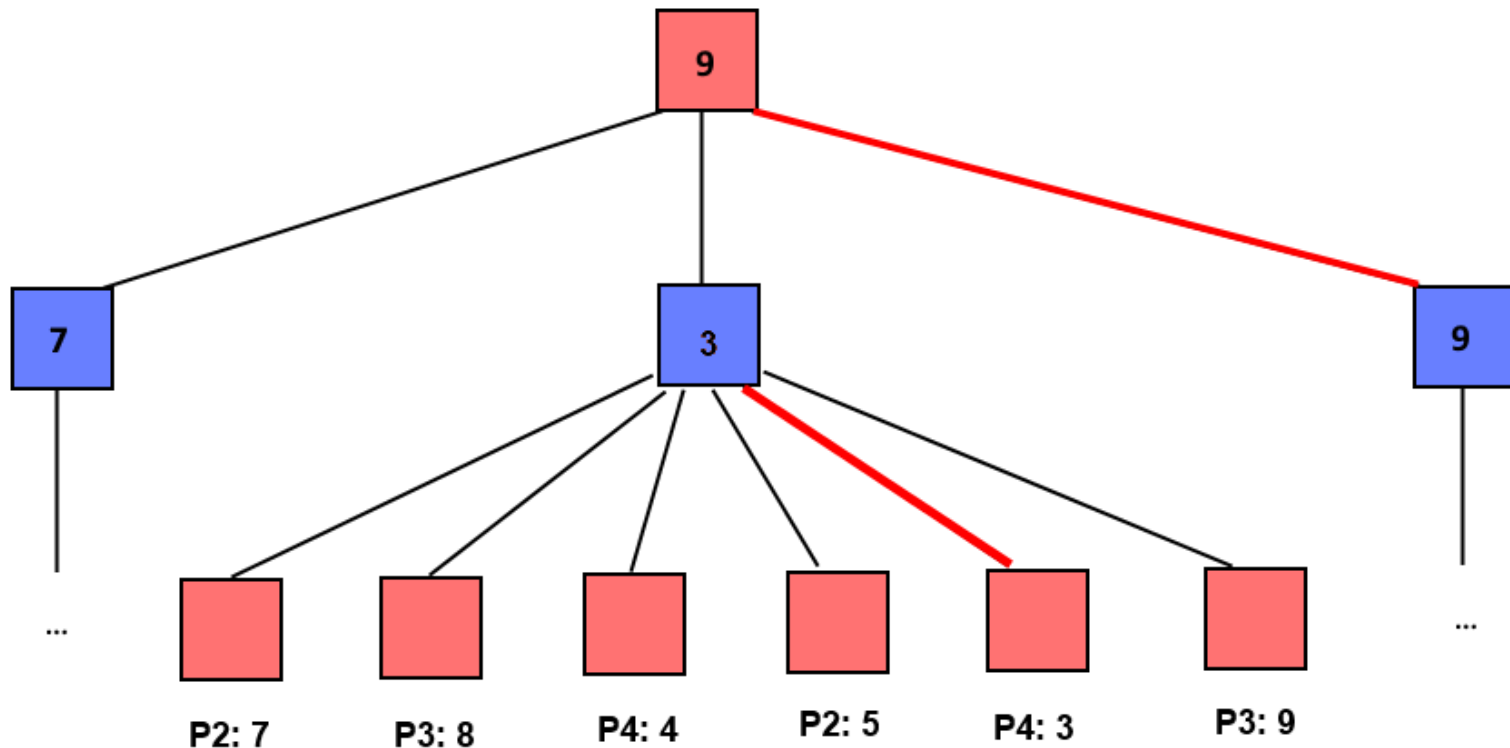
Max-N Algorithm

- In terms of raw Mini-Max, very simple extension
- Pros
 - Players “look out for number one”
 - More realistic play
 - Perspective player can see more opportunities
 - Reason: Possibilities are not excluded as readily
- Cons
 - Pruning is very complicated, and not as good
 - Can be worse than Paranoid due to decreased search depth

Best-Reply Search

- Relatively new: 2011 (Schadd and Winands)
- *All* opponents considered to be *one* player
 - They only get ONE turn between them
- Only opponent with best move is thought to act
- Return to MAX-MIN-MAX-MIN...
- Essentially a return to Mini-Max algorithm
 - *With a very powerful* opponent!!

Best-Reply Search



Sample BRS Tree.

The children of the other minimizing nodes are omitted.

Best-Reply Search

```
function integer best-reply(node, depth):
  if node is terminal or depth <= 0 then
    return heuristic value of node
  else
    if node is max then
      val =  $-\infty$ 
      for all child of node do
        val = max(val, best-reply(child; depth - 1)
      end for
    else
      val =  $\infty$ 
      for all opponents do
        for all opponent's child at node do
          val = min(val; best-reply(child;
            depth - 1)
        end for
      end for
    end if
  end if
end if
```

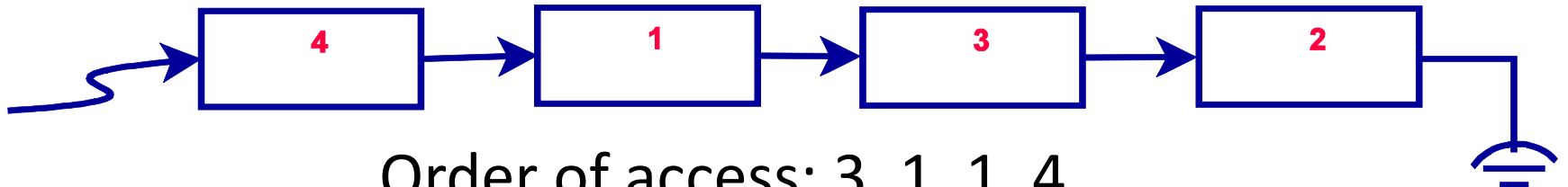
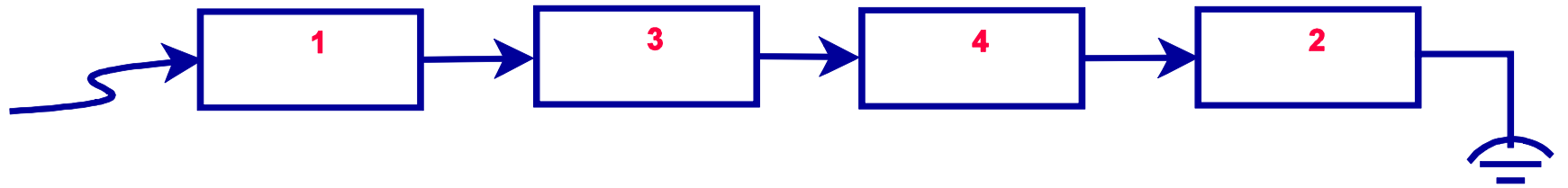
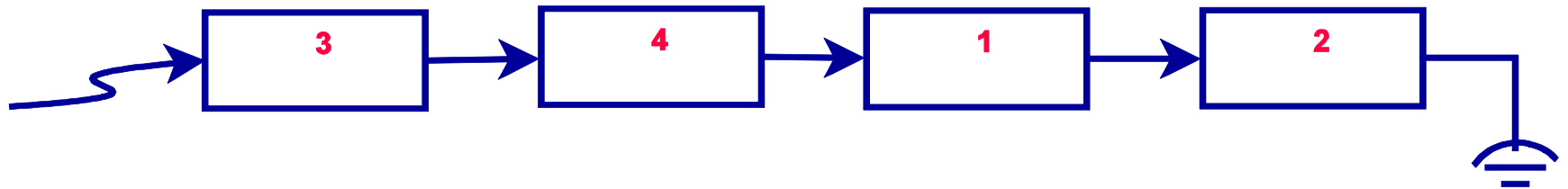
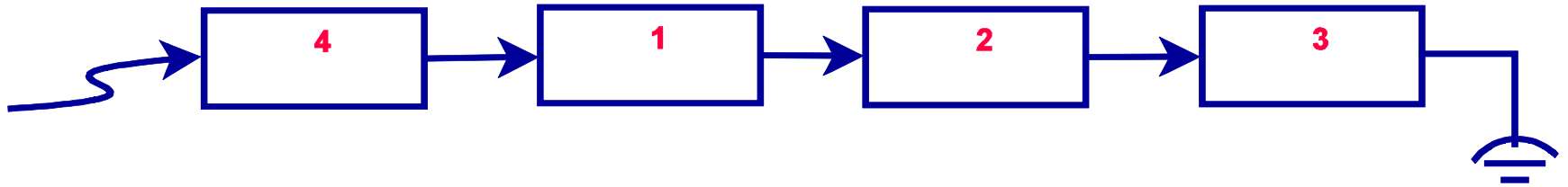
Best-Reply Search

- Attempt to get “best of both worlds”
- Pros
 - Balance between coalition and free-for-all
 - Allows α - β pruning
 - Significant lookahead for perspective player
- Cons
 - Illegal game states analyzed
 - Not applicable to some games
 - This is the domain of some current research (2015)

Adaptive Data Structures

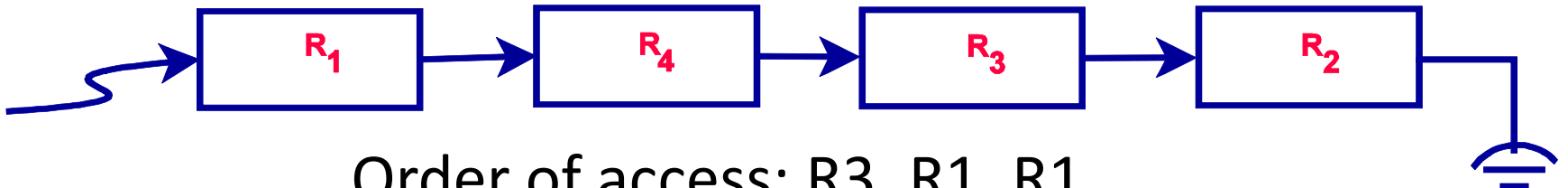
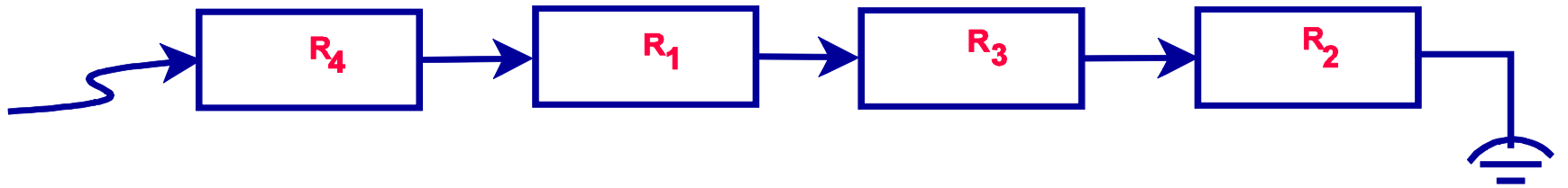
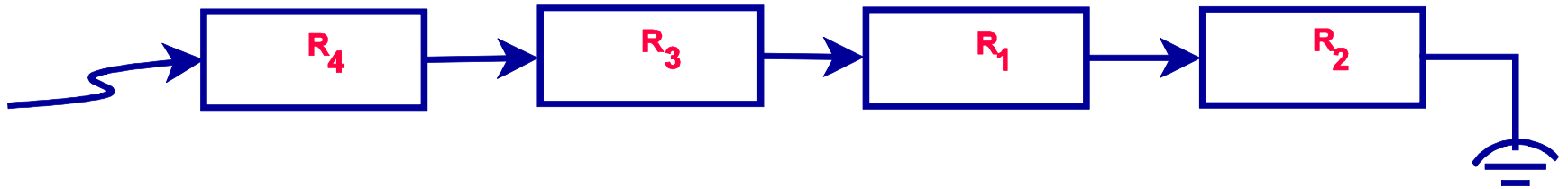
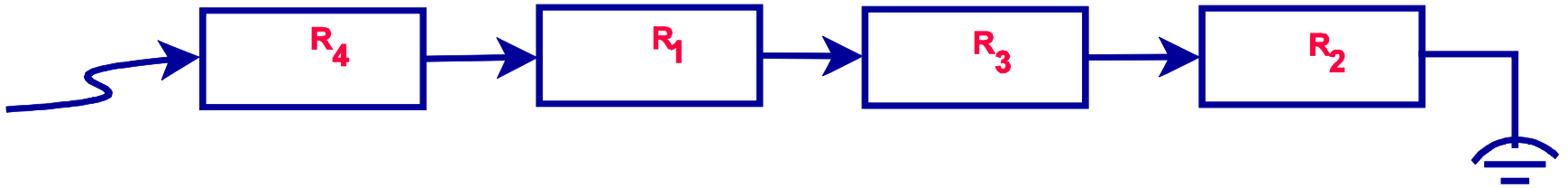
- Other, completely unrelated field
- Concerned with record access frequency
- Problem:
 - Elements in data structure accessed with different frequency
- Solution:
 - Change the structure of the data structure as elements queried
- Can use list, tree or others

ADS – Move to Front Rule



Order of access: 3, 1, 1, 4, ...

ADS – Transposition Rule

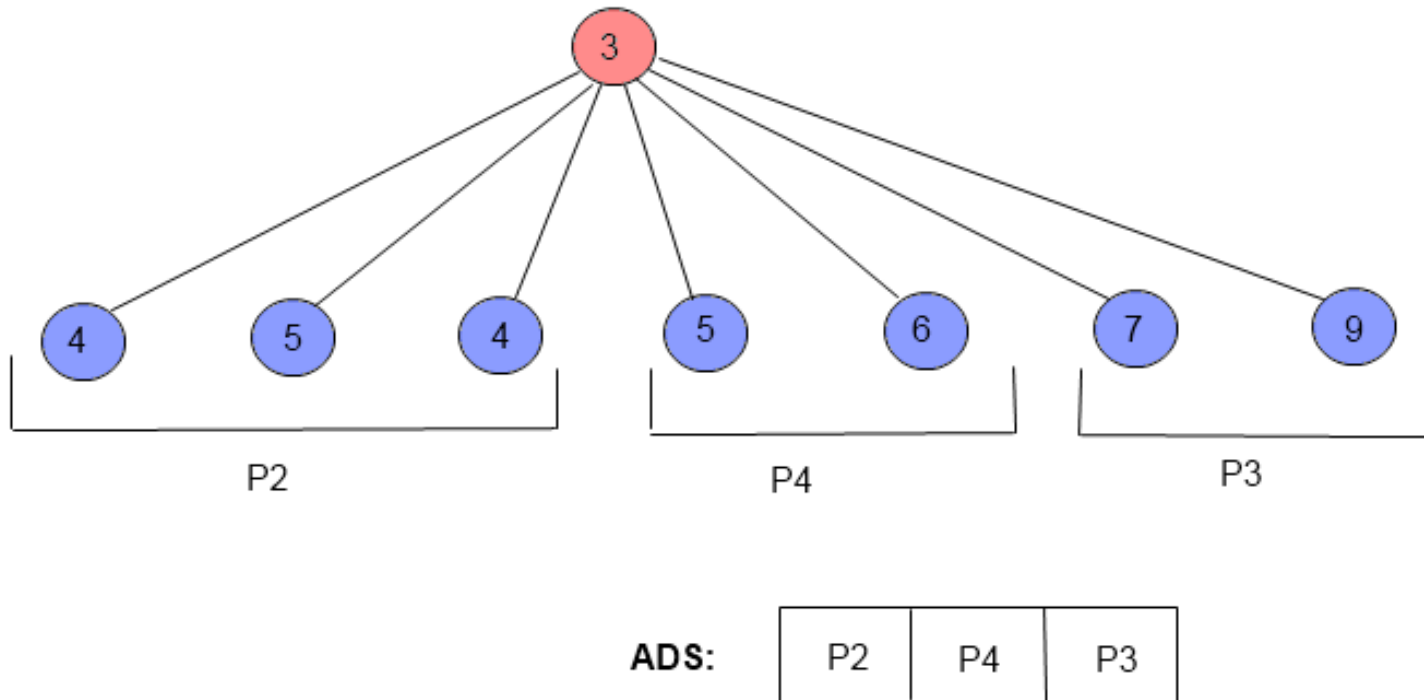


Order of access: R_3, R_1, R_1, \dots

The Threat-ADS Heuristic

- Our contribution, usable with the BRS
- ADS operations are constant, and small
- We use an ADS that contains **opponents**
- When an opponent is found to have the most *minimizing move*, we query the ADS
- ADS moves over time to **relative opponent threats**
- When grouping moves, do it in the order of the ADS
- Improves **move ordering**, leading to better pruning!

The Threat-ADS Heuristic



BRS with Threat-ADS (one level)

BRS with Threat-ADS

```
function integer brs_threat_ads(node, depth):
  if node is terminal or depth <= 0 then
    return heuristic value of node
  else
    if node is max then
      val =  $-\infty$ 
      for all child of node do
        val = max(val, best-reply(child; depth - 1)
      end for
    else
      val =  $\infty$ 
      for all opponents in ADS do
        for all opponent's child at node do
          val = min(val; best-reply(child; depth - 1)
        end for
      end for
      ADS.update(val.opponent)
    end if
  end if
```

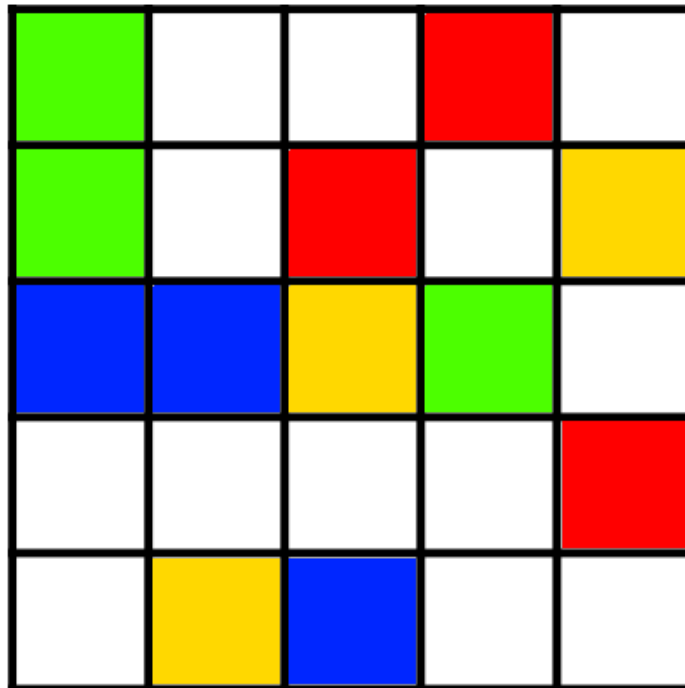

Experimental Framework

- Game needed to test Threat-ADS heuristic
- Needs:
 - BRS must be applicable
 - Game should be simple to implement
- Use established games Focus and Chinese Checkers
- Also develop the Virus Game

Virus Game

- Turn based game with N players
- Played on 2D board
- Goal is to eliminate all other players
- Turn: Player “infects” a square they are adjacent to
- All nearby squares, according to a configured pattern, are given to that player

Virus Game



Experimental Configuration

- One player: BRS with Threat-ADS
- Others: Random (Interested in *tree pruning*)
- Take *Node Count* over first few turns of the game
 - Count each node expanded, but not those pruned
- Average over 50 games
- Run for each of three games mentioned
- Run over a variety of configurations
 - Varying number of players
 - Varying starting state

Results (Initial Board State)

Game	Threat-ADS?	Avg. Node Count
Virus Game	No	264,000
Virus Game	Yes	237,000
Focus	No	6,859,000
Focus	Yes	6,443,000
Chinese Checkers	No	3,485,000
Chinese Checkers	Yes	3,070,000

Results (Midgame Board State)

Game	Threat-ADS?	Avg. Node Count
Virus Game	No	307,000
Virus Game	Yes	275,000
Focus	No	14,460,000
Focus	Yes	13,050,000
Chinese Checkers	No	8,170,000
Chinese Checkers	Yes	7,680,000

Monte-Carlo Methods

- *Entirely* different way of looking at game playing
 - Applicable to two player and multi-player games
- No game heuristics required!
- Driven by *random game playing*
 - Strong when no good heuristic is available
 - Big example in research is Go
- Very simple example:
 - Play 50 random games for each move
 - Pick one with highest winrate

Monte-Carlo Tree Search

- Simple example above
 - Works for easy games
 - Look-ahead is useful
- Apply random game playing to game tree search
- Navigate:
 - From root to unvisited node
 - Then play random game(s)
- Path guided by exploration/exploitation balance
- At end of time, pick most promising move
- Very powerful: Relatively new compared to Mini-Max

UCT Algorithm

- Dominant Monte-Carlo Tree Search technique (2015)
- Starting from root:
 - If there is an unvisited child, pick it
 - Otherwise, pick child that maximizes **UCTValue**
$$\text{UCTVal} = \text{winrate} + \text{sqrt}(\ln(\text{parent.visits})/\text{visits})$$
 - Repeat until an unvisited child is found
- Propagate winrate back up to root
- Repeat until time is up
- Pick move that has *highest winrate*

UCT Algorithm

```
function integer uct(node, depth):  
  for time-steps do  
    position = root  
    while position is explored  
      val = -∞  
      for child of position  
        !– Unexplored node check here--!  
        val = max(val, UCTValue(child))  
        position = val.node  
      end for  
    end while  
  
    Play random game(s) at child  
    while position is not root  
      update win-rate for player at node  
      position = position.parent  
    end while  
  
  end for
```

Multi-Player UCT Algorithm

- *Very easy* to extend
 - We do not have to maintain heuristic values
 - UCT handles N-player games in its base form
- For the player making the move
 - Simply record winrate at each node
 - Assume player will pick move most likely to lead to win
- No change from previous algorithm

More on UCT Value

- UCTValue has two parts
- Winrate is self-explanatory
 - Value between 0.0 and 1.0 indicating proportion of wins
- Second part: **$\sqrt{\ln(\text{parent.visits})/\text{visits}}$**
- Specifics not important, but also between 0.0 and 1.0
- Goes *up* the less this child has been explored in relation to its parent
- Achieves exploration/exploitation balance!
- Sometimes constants usually added to tweak this

Applications of UCT

- Best performance available for Go
 - Top player is currently Zen
 - Defeated 9-dan player with three stone handicap
- Applied to wide range of games
 - Poker
 - Settlers of Catan
 - Magic: The Gathering